



A Consensus Mechanism Based on an Improved Genetic Algorithm

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Abstract

An important feature of blockchain technology is that all participants jointly maintain transaction data and can achieve mutual trust relationships without integrated control, which relies on distributed consensus algorithms. Practical Byzantine Fault Tolerant algorithm (PBFT) is a fault-tolerant algorithm based on state machine replication, which solves the Byzantine error, that is, the malicious behavior of nodes. In PBFT, all participating nodes are divided into the primary node and backup nodes. When this primary node commits evil or fails, it will elect a primary node again for message communication. The genetic algorithm (GA) is a computer simulation study inspired by the natural biological genetic evolution criterion “natural selection, survival of the fittest”. Genetic algorithm is actually a method to find the optimal solution. According to it, the best primary node is selected in the PBFT algorithm to improve consensus efficiency. The consensus algorithm is the guarantee of the decentralization feature in blockchain technology. The PBFT algorithm is a commonly used consensus algorithm. However, this algorithm has the following problems: when the primary node fails, it must be selected again, which leads to a decrease in consensus efficiency. This paper proposes a consensus mechanism based on an improved genetic algorithm, which uses an improved genetic algorithm to select the primary node. According to the genetic algorithm, the best primary node is selected, and it meets the minimum number of errors or evils and the highest transaction efficiency with other backup nodes. The improved consensus algorithm can effectively reduce system delay and improve consensus efficiency.

Subject Areas

Bioinformatics

Keywords

Consensus Mechanism, Genetic Algorithm, Improvement

1. Introduction

In blockchain system, the consensus algorithm enables all participants to jointly maintain transaction data and achieve mutual trust relationships without integrated control. Consensus algorithms can be roughly divided into proof-based consensus algorithms and voting-based consensus algorithms; the most classic consensus voting-based algorithm is the PBFT algorithm. But this consensus algorithm has some problems to improve. Then many variant consensus algorithms are presented to issue the shortcomings of PBFT. In BBFT [1], it uses the topology tree structure and message aggregation to reduce the communication complexity. Liu *et al.* presented the FastBFT [2] in 2018. It uses a message aggregation and trust execution environments technique to improve the PBFT. Since in PBFT, the total number of nodes is $3f + 1$ (f is the number of fault nodes), then Giuliana Santos Veronese *et al.* propose MinBFT [3] and Rudiger Kaitza *et al.* present CheapBFT [4]. Their methods all reduce the total nodes' number from $3f + 1$ to $2f + 1$. Zyzzyva [5] is a Speculative Byzantine Fault Tolerance which is proposed by Ramakrishna Kotla *et al.* It changes three-protocol of PBFT to improve performance. HAO *et al.* proposed the CDBFT [6]. It combines DPOS and PBFT to reduce the participation probability of malicious replicas in the consensus process. Jeon *et al.* proposed the DBFT [7] which makes optimization for the application of public blockchain. Among all the PBFT variant algorithms, few have improved the algorithm by changing the way the master node is selected. So this article proposes a consensus mechanism based on an improved genetic algorithm. BAGLEY took the lead in using the concept of "genetic algorithm" in the paper [8] in 1967. The improved genetic algorithm is used to obtain the best primary node when it is selected. After calculation, the algorithm improves transaction efficiency and processing time. The structure of the paper is as follows: Section 2 describes the traditional PBFT algorithm, Section 3 introduces the consensus algorithm based on the improved genetic algorithm, and Section 4 does a summary to this article.

2. Practical Byzantine Fault Tolerance Algorithm

PBFT algorithm was presented by Castro and Liskvo in 1999 [9]. It requires that each consensus be executed in the same order on each copy, and each participant decides the content of the consensus by voting. PBFT adopts a three-phase protocol: PREPREPARE, PREPARE, and COMMIT.

PRE-PREPARE

In the PRE-PREPARE phase, the primary sends a pre-prepare message to all replica nodes.

PREPARE

When the backup node i has received the pre-prepare message, it enters the PREPARE phase, and sends the prepare message to other nodes. Besides it writes both the pre-prepare message and the prepare message to the log.

In the PBFT algorithm, the PRE-PREPARE phase and the PREPARE phase

ensure that the non-faulty replica nodes agree on the ordering of requests in the same view. BFT systems shall guarantee safety and liveness in the presence of faulty servers and clients [10].

COMMIT

The replica joins in the COMMIT phase and broadcasts to all replicas the commit message. **Figure 1** is a flow chart of the three-phase agreement.

The above three steps determine the number of replicas in PBFT is $3f + 1$ and the communication complexity is $O(f^2)$ (f is the number of faulty nodes).

3. A Variant Consensus Mechanism Based on an Improved Genetic Algorithm

Since in the PBFT algorithm nodes are divided into the primary node and other replica nodes, and the primary node as a consensus node executes a three-phase protocol to complete transactions. Once the primary node fails, it will re-select, which will cause the efficiency of system transaction execution to decrease. Therefore, the selection of the primary node is very essential. This paper proposes an improved genetic algorithm to select the best primary node of PBFT. The improved genetic algorithm mainly uses the elitist strategy to retain the best individual, and the best individual does not undergo mutation, crossover and other operations, and directly enters the next generation. The improved algorithm not only improves convergence, but also ensures that the optimal individual characteristics are not destroyed in the evolution process [11] [12]. Firstly, each participant is involved in coding into the corresponding chromosomes to form the initial initial population. Secondly, select the master node according to the fitness function. Then use the elitist selection strategy to perform crossover and mutation operations to achieve the best primary node selection.

Step I

Encoding all participating codes to chromosome. Suppose there are m participating nodes, a chromosome of an m -dimensional vector is obtained after encoding, this dimension is the number of all participating nodes $N = \{n_1, n_2, \dots, n_p, \dots, n_m\}$.

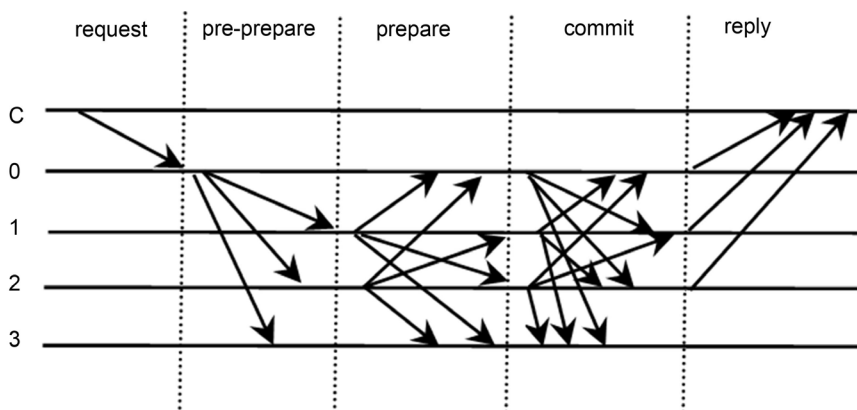


Figure 1. PBFT algorithm flow chart.

Step II

Setting fitness function $f(n_i) = \sum_{n_i \in N} \text{Fau}(n_i) \text{EFF}(n_i)$ for all encoded chromosomes. n_i represents a legal chromosome, N represents the set of all chromosomes, $\text{Fau}()$ represents the number of times that the node has problems, and $\text{Eff}()$ represents the efficiency of the node's transaction with other nodes. According to the $f(n_i)$ to obtain the optimal solution, the best primary node is selected.

Step III

Selecting the best master node according to the elitist selection strategy, and other nodes perform mutation and selection operations.

Step IV

When the termination conditions are met, the optimal solution of the function will be used as the primary node to participate in the implementation of the PBFT three-phase protocol.

Figure 2 is a flowchart of the improved algorithm.

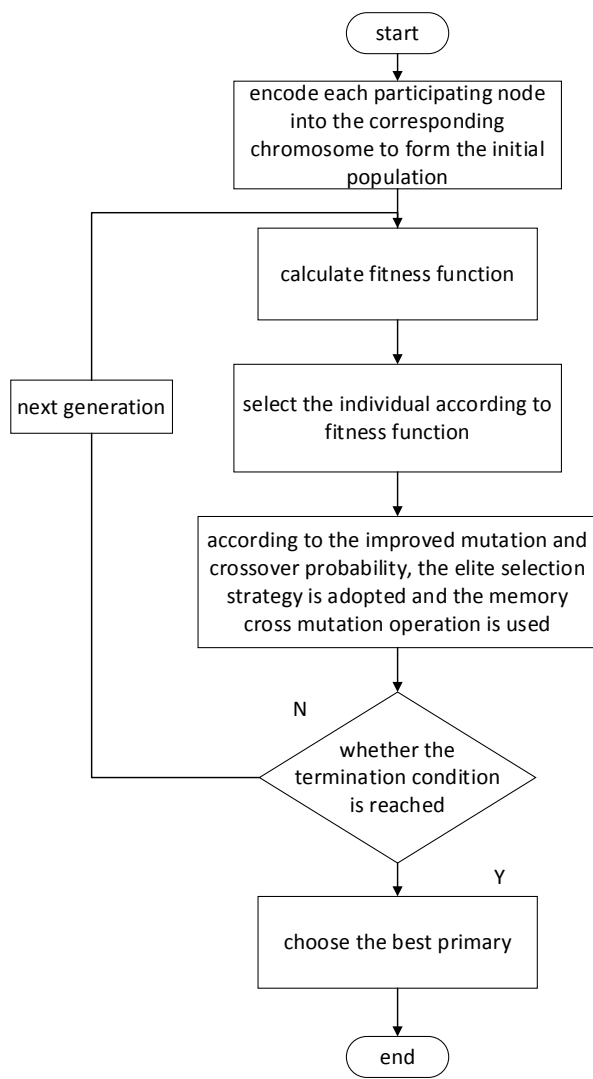


Figure 2. Improved algorithm flow chart.

4. Conclusions

Latency is an important indicator to measure the performance of the blockchain. The delay is mainly affected by two aspects: one is the network communication performance when the message is propagated between nodes, and the other is the operating efficiency of the consensus algorithm. According to formula $\text{Delay} = T_{\text{Broadcast}} + T_{\text{Consensus}}$, in the fitness function $f(n_i) = \sum_{n_i \in N} \text{Fau}(n_i) \text{EFF}(n_i)$, if a node meets the minimum $\text{Fau}()$ and the maximum $\text{Eff}()$, during the consensus process, $T_{\text{Broadcast}}$ decreases, so the entire Delay decreases. Compared with the traditional PBFT algorithm, the delay is effectively reduced.

This paper proposes a PBFT algorithm based on an improved genetic algorithm, which uses genetic algorithm to improve the selection strategy of the primary node, so that the improved PBFT algorithm consensus efficiency is higher.

Conflicts of Interest

The authors declare no conflicts of interest regarding the publication of this paper.

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