

Equivalence between Linear Tangle and Single Ideal

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Abstract

The concept of linear tangle was introduced as an obstruction to mixed searching number. The concept of single ideal has been introduced as an obstruction to linear-width. Moreover, it was already known that mixed search number is equivalent to linear-width. Hence, by combining those results, we obtain a proof of the equivalence between linear tangle and single ideal. This short report gives an alternative proof of the equivalence.

Keywords

Linear Tangle, Single Ideal, Submodular Function

1. Introduction

A graph searching game is a game where searchers (or cops) want to capture a fugitive (or robber) and the fugitive want to escape from the searchers, and they move through a graph for their purpose. Graph searching games have been well-studied [1] [2] [3], since graph searching games have many practical and theoretical applications such as robot motion planning, network security, and artificial intelligence (see e.g. [4]).

There are several variants of graph searching games such as edge search, node search, and mixed search (see e.g. [5]) and there are several graph width-parameters such as path-width, tree-width, and branch-width. In the study of graph searching games, it is known that there are some strong connections between graph searching games and graph width parameters, in which the minimum number of searchers (*i.e.*, *search number*) usually corresponds to the value of width. For example, mixed search number is essentially equivalent to *linear-width*¹ (see [6] [7]).

¹If a graph has no pendant vertices, the equivalence holds.

The concept of linear tangle was introduced in [7] as an obstruction to the existence of mixed searching strategy: there is a linear tangle of order $k+1$ iff there is no mixed searching strategy with k searchers, where k is a prefixed integer. From the equivalence, as mentioned above, between mixed search number and linear-width, a linear tangle of order $k+1$ is an obstruction to being linear-width is at most k (see also [8]).

The concept of single ideal has been introduced in [9] as an obstruction to linear-width: there is a single ideal of order $k+1$ iff the linear-width is more than k . Thus, a linear tangle of a large order and a single ideal of a large order are both obstructions to small linear-width, which means that the concepts of linear tangle and single ideal are the same. In this short report, we give an alternative proof of the equivalence between linear tangle and single ideal.

2. Definitions and Notations

In this paper, we consider a pair (E, f) rather than graphs, where E is an underlying set and f is a symmetric submodular function on E , and such a pair is called *connectivity system* (see cf. [10]). All sets considered in this paper are finite. For an underlying set E and a subset X of E , we denote $E \setminus X$ by \bar{X} .

A function $f : 2^E \rightarrow \mathbb{Z}$ is *symmetric submodular* if f satisfies the following:

- 1) $f(X) = f(\bar{X})$ for any $X \subseteq E$,
- 2) $f(X) + f(Y) \geq f(X \cup Y) + f(Y \cap X)$ for any $X, Y \subseteq E$.

It is known that a symmetric submodular function f satisfies the following properties [10]: for each $X, Y \subseteq E$,

- 1) $f(X) \geq f(\emptyset)$ and
- 2) $f(X) + f(Y) \geq f(X \setminus Y) + f(Y \setminus X)$.

A set X is *k-efficient* if $f(X) \leq k$. Throughout the paper, f means a symmetric submodular function, k a fixed positive integer, and we assume that $f(\{e\}) \leq k$ for every $e \in E$, hence we have $f(\emptyset) \leq k$.

Definition 1 ([7]). A *linear tangle* of order $k+1^2$ on a connectivity system (E, f) is a family \mathcal{L} of k -efficient subsets of E , satisfying the following axioms:

- (L1) $\emptyset \in \mathcal{L}$,
- (L2) For each k -efficient subset X of E , exactly one of X or \bar{X} in \mathcal{L} ,
- (L3) If $X, Y \in \mathcal{L}$, $e \in E$, and $f(\{e\}) \leq k$, then $X \cup Y \cup \{e\} \neq E$ holds.

Definition 2 ([9]). A *single ideal* of order $k+1$ on a connectivity system (E, f) is a family \mathcal{M} of k -efficient subsets of E , satisfying the following axioms:

- (S1) $E \notin \mathcal{M}$,
- (S2) If $A, B \subseteq E$, $A \subset B$, $B \in \mathcal{M}$, and $f(A) \leq k$, then $A \in \mathcal{M}$ holds,
- (S3) If $A \in \mathcal{M}$, $e \in E$, $f(\{e\}) \leq k$, and $f(A \cup \{e\}) \leq k$, then $A \cup \{e\} \in \mathcal{M}$ holds.

We also consider the following additional axiom:

²In [7], the order is k rather than $k+1$.

(S4) For each k -efficient subset A of E , exactly one of A or \bar{A} in \mathcal{M}

It is shown in [9] that the linear-width of (E, f) is at least $k+1$ if and only if there is a single ideal on (E, f) of order $k+1$ which satisfies (S4).

Since we assume that $f(\{e\}) \leq k$ for every $e \in E$, the axioms (L3) and (S3) can be restated, respectively, as follows:

(L3) If $X, Y \in \mathcal{L}$ and $e \in E$, then $X \cup Y \cup \{e\} \neq E$ holds.

(S3) If $A \in \mathcal{M}$, $e \in E$, and $f(A \cup \{e\}) \leq k$, then $A \cup \{e\} \in \mathcal{M}$ holds.

3. Result

Lemma 1. A linear tangle \mathcal{L} of order $k+1$ is a single ideal of order $k+1$ satisfying the additional axiom (S4).

Proof. From the axioms (L1) and (L2), it is obvious that \mathcal{L} satisfies the axioms (S1) and (S4).

We claim that \mathcal{L} satisfies the axioms (S2). Suppose, to the contrary, that there exist k -efficient subsets A and B such that $A \subseteq B$, $B \in \mathcal{L}$, and $A \notin \mathcal{L}$. Then, we have $\bar{A} \in \mathcal{L}$ by (L2), and for any $e \in E$, $\bar{A} \cup \{e\} \cup B = E$ holds, but this contradicts the axiom (L3).

Finally, we show that \mathcal{L} satisfies the axioms (S3). Suppose, to the contrary, that there exists k -efficient subset $A \in \mathcal{L}$ and an element $e \in E$ such that $f(A \cup \{e\}) \leq k$ and $A \cup \{e\} \notin \mathcal{L}$ hold. Then, we have $\overline{A \cup \{e\}} \in \mathcal{L}$, hence $A, A \cup \{e\} \in \mathcal{L}$ and $A \cup A \cup \{e\} \cup \{e\} = E$ hold, however, this contradicts the axiom (L3).

Lemma 2. A single ideal \mathcal{M} of order $k+1$ satisfying the additional axiom (S4) is a linear tangle of order $k+1$.

Proof. From the axioms (S1) and (S4), it is obvious that \mathcal{M} satisfies the axioms (L1) and (L2).

We show that \mathcal{M} satisfies the axiom (L3). Suppose, to the contrary, that there exists a triple $(X, Y, \{e\})$ such that $X, Y \in \mathcal{M}$, $e \in E$, and $X \cup Y \cup \{e\} = E$. We choose a triple minimizing $|X \cap Y|$ in such triples $(X, Y, \{e\})$. First, we claim that $X \cap Y = \emptyset$ holds. Since $2k \geq f(X) + f(Y) \geq f(X \setminus Y) + f(Y \setminus X)$ holds, at least one of $f(X \setminus Y)$ or $f(Y \setminus X)$ is at most k . Without loss of generality, we may assume that $f(X \setminus Y)$ is at most k . Hence, by (S2) from $X \setminus Y \subseteq X$, we have $X \setminus Y \in \mathcal{M}$. If $X \cap Y \neq \emptyset$, then we have $|X \cap Y| > |(X \setminus Y) \cap Y|$, however, this contradicts the choice of the triple. Thus, we have shown that $X \cap Y = \emptyset$.

Next, we claim that $e \notin X$ holds. Suppose if not, that is, if $e \in X$, then we have $X \cup Y = E$, which implies that $\bar{X} = Y \in \mathcal{M}$ holds, however, this contradicts the axiom (S4). Similarly, we know that $e \notin Y$ holds.

Now, we know that the triple $(X, Y, \{e\})$ consists of a partition of E . Hence, we have $f(X \cup \{e\}) = f(\bar{Y}) = f(Y) \leq k$, and from this, it follows that $X \cup \{e\} = \bar{Y} \in \mathcal{M}$ holds by the axiom (S3). However, this contradicts the axiom (S4).

From lemmas 1 and 2, we have the following theorem.

Theorem 1. *Under the assumption that $f(\{e\}) \leq k$ for every $e \in E$, \mathcal{F} is a linear tangle of order $k+1$ iff \mathcal{F} is a single ideal of order $k+1$ satisfying the additional axiom (S4).*

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Conflicts of Interest

The authors declare no conflicts of interest regarding the publication of this paper.

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