

# Agent-Based Synthesis of Distributed Controllers for Discrete Manufacturing Systems

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## ABSTRACT

*A method for designing real-time distributed controllers of discrete manufacturing systems is presented. The approach held is agent based; the controller strategy is distributed into several interacting agents that operate each one on a part of the manufacturing process; these agents may be distributed into several interconnected processors. The proposed method consists of a modelling methodology and software development framework that provides a generic agent architecture and communication facilities supporting the interaction among agents.*

**Keywords:** *Agent-Based Software Design, Distributed Knowledge Based Controllers, Discrete Manufacturing Systems*

## 1. Introduction

Nowadays discrete manufacturing systems are large and complex systems that integrate several kinds of devices of miscellaneous nature and behaviour, namely robots, conveyors, machines, sensors, etc. Additionally, the production requirements are often changed; these facts impose to the system components to be versatile, and to the coordination system or controller to be highly flexible.

The core of a control system is complex software, which determines at last, the flexibility and the performance of the automated system. This control software provides several functions: tasks execution, monitoring, decision making, and planning; it is generally distributed into a four layer hierarchy [1]; in this scheme the control function is decomposed into four levels in which the response time is shortest in the lower levels.

The lowest level includes the local controllers of the physical devices in the cell (robots, conveyors, machines, sensors, etc). The task coordination level or cell level manages and supervises the activities of the local controllers involved in a cell by the generation of pertinent commands according to events issued from the local control level. The task planner generates the strategy of the controller for the cells according to the specifications from the production planning level.

Due to the complexity of the tasks, especially which performed at the cell level, the synthesis of the controller is a difficult job often addressed through planning tech-

niques. One of the problems found in the synthesis real-time controllers is that the representation of the qualitative controller often includes a large amount of knowledge whose processing is time consuming.

This work deals with the task coordinator level. The functions of this control level can be decomposed mainly in a) the sequencing of operations to accomplish the assembly task in a normal functioning regime, and b) the handling of exceptions representing operation failures [2]. In this paper the case of normal execution is addressed.

The design and implementation of tasks controllers of complex manufacturing systems has been addressed by using several approaches. The object oriented approach has been held for modelling [3], simulation [4-6], and control [7] of manufacturing systems. The agent based approach [8] has been adopted to address some problems in manufacturing systems [9]; DeLoach [10] proposes a modelling language for describing the diverse kinds of agents, and defines a methodology (MaSE) for the formal synthesis of agent systems; in [11] Bussmann focuses on decision making issues during the planning stage, and in [12] he addresses the task programming issue proposing a synthesis method that leads to concurrent centralised control software; in [13] Ouelhadj proposed a dynamic control architecture for manufacturing systems organised into cells, but the programming of the agents is not reported.

In this work we also profit of the agent-based approach for conceiving a control software as a composition of interacting modules, defined as reactive agents, which may

be distributed into several processors; it is proposed a method that supports the complete development life cycle of distributed controllers of discrete manufacturing systems; this lifecycle is shown in **Figure 1**, in which the stages of the method are pictorially overviewed. The proposed method for sequencing the activities of the cell components allows building rapidly prototypes of software controllers.

The remainder of this paper is organised as follows: Section 2 describes the proposed methodology for modelling both the manufacturing system and the tasks to be executed; Section 3 presents the proposed method for the design of distributed control software: first the decomposition of the task model is described, then agent based solution to the synthesis of distributed software is outlined.

## 2. Controller Modelling

This section presents the methodology that helps to obtain systematically the contents of the knowledge bases from a model of the manufacturing/assembly system; this model, close to that presented in [7], includes the system description and the tasks specification. The aim of this stage is to obtain in a structured way the system functioning and production requirements.

The description consists in a component classification of the manufacturing process, and a structuring of the system workspace.

The component classification leads to a taxonomy of the system components organised as a hierarchy including capabilities and features of each component, and the total quantity of devices. The hierarchy is useful to program the necessary classes in the control software. **Figure 2** shows an example of components hierarchy.

### 2.1. System Description

The structuring of the system workspace consist of a definition of relevant physical emplacements where operations are performed on the work pieces or parts; rather than space partition into regions, the structuring is a discrete assignment of positions. The key element for structuring the workspace is the notion of *site*, which is defined as the place where parts can be temporary held or stored in a stable position (a table, a magazine, a robot gripper, ...) [2].

The sites may be single or composed (macro-site); single sites held one part or subassembly; composed sites have two or more emplacements which manage the information attached to a set of sites closely located and functionally equivalents. The sites that are associated to effectors are named *active sites*; otherwise they are called *passive sites*. A site contains information about the work zone were it is emplaced that can be used as mutual exclusion resource (for robot collision avoidance, for example).

## 2.2. Task Specification

The flow of material is described by a flow of parts graph (FPG), and then a set of dispatching rules, which represent the controller strategy, is obtained.

### 2.2.1. The Flow of Parts Graph

A FPG is a directed graph whose nodes are all the sites of the system; the arcs joining the nodes represent either the operations needed to transfer the parts from one site to another one or to modify the properties of the part held into a site. For example, **Figure 3** describes the operations *pick* and *place* performed by a robot (R1); three sites are involved: two passive sites (CONV2 and TAB1), and an active site (GRIP1) associated to the robot gripper. The definition of FPG is given below:

**Definition.** A flow of parts graph is the tuple  $F = (G, \text{SITES}, \text{OPER}, \varepsilon, \lambda, \phi)$ , where

- $G$  is a connected directed graph  $G = (V, A)$ , where
  - $V$  is a finite set of vertex,
  - $A \subseteq V \times V$  is a set of edges or arcs.
- $\text{SITES} = \{\text{site}_1, \dots, \text{site}_n\}$  is a finite set of site names, which are not input or output sites.
- $\varepsilon = \{\text{sitein}_1, \dots, \text{sitein}_p, \text{siteout}_1, \dots, \text{siteout}_q\}$  is a finite set of site names labelling sites where the parts entry or leave the FMS.
- $\lambda: V \rightarrow \text{SITES} \cup \{\varepsilon\}$  is a labelling function that assigns name sites to the vertex of  $G$ .
- $\phi: A \rightarrow \text{OPER}$  is a labelling function that assigns operations names to the arcs of  $G$ .

### 2.2.2. Sequencing the Operations

The dispatching rules are antecedent-consequent rules that state the conditions in which an operation must be executed; they are obtained directly from the FPG, and the number of rules is the same than the number of arcs in the FPG. The antecedent part is composed by conditions that involve mainly sensory conditions and tests (contents, part posture, ...) on the sites related by the operation; other conditions may involve tests on sites located upstream the FPG. The consequent part includes the request of execution of the associated operation and the updating of the involved sites.

## 2.3. A Modelling Example

For illustrating purposes we include an example regarding a simple assembly system; it will be addressed through the rest of the paper.

### 2.3.1. System and Task Description

1) **The system:** Consider the assembly cell sketched in **Figure 4**; it consist of three conveyor belts B1, B2, and B3, two robots R1, R2, two assembly tables A1, A2, an storing table ST. Each assembly table has two positions;

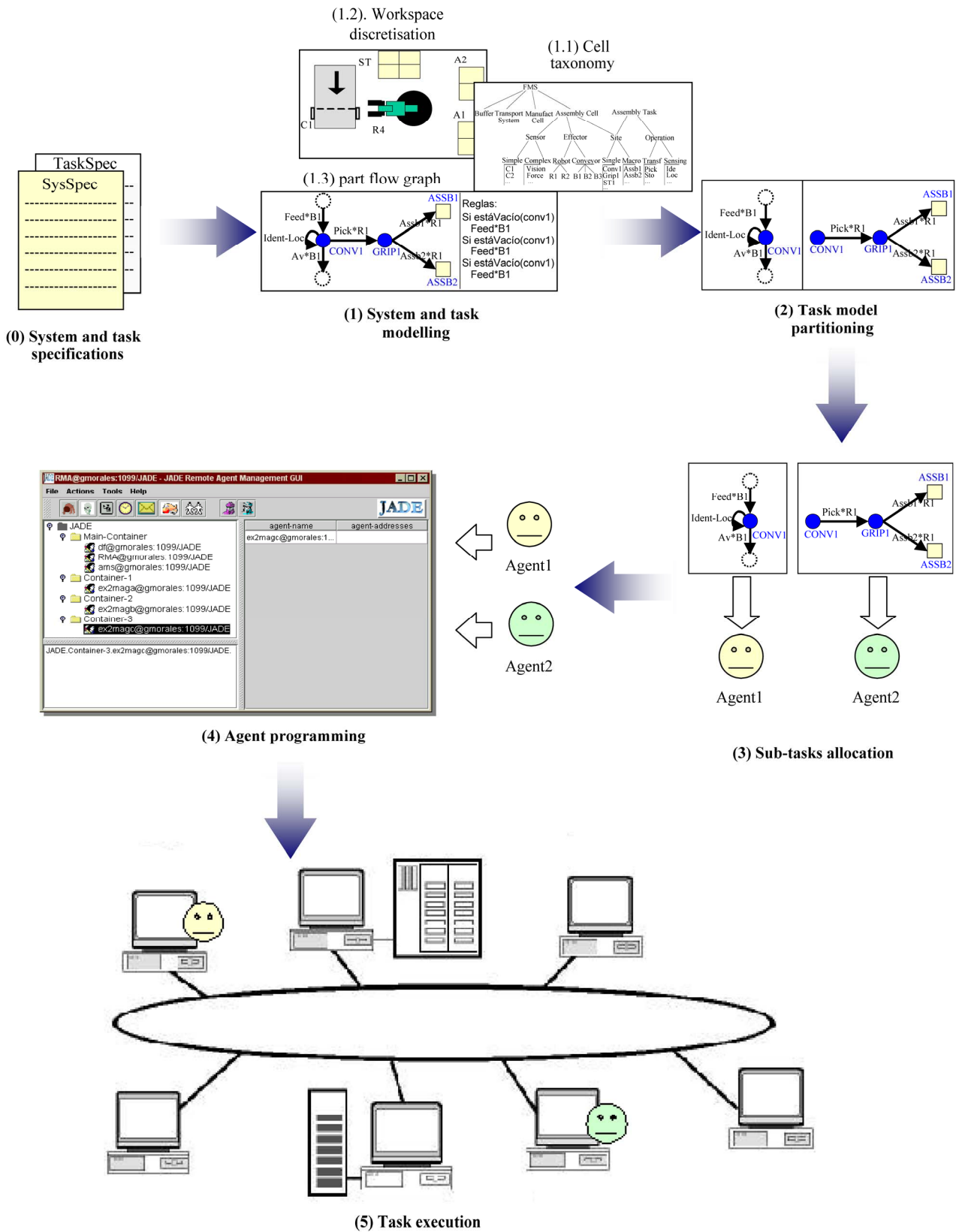


Figure 1. Software development life cycle.

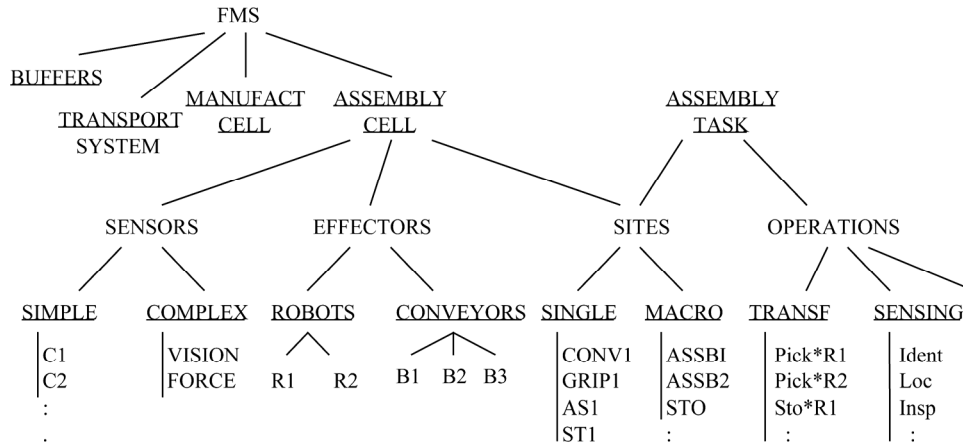


Figure 2. Taxonomy of the assembly cell.

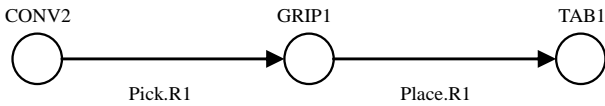


Figure 3. Flow of parts graph.

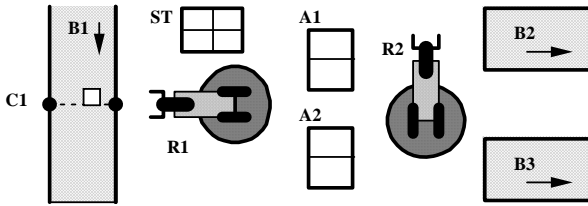


Figure 4. Assembly cell layout.

the storing table has four positions. The parts to be assembled arrive through the conveyor B1; two kinds of assembled products leave the cell through B2 and B3. In the front of R1, an optical sensor C1 detects the arrival of parts; over this zone a camera of a location-recognition system is emplaced.

2) **The task:** Eight types of parts (A, B, C, D, E, F, G, and H) constitute the input flow in B1; they arrive at random order. R1 gets the parts and builds assemblies (stacks) on A1 (with the parts A, B, C, and D) or A2 (with the parts E, F, G, and H) according to a predefined order for each product. The detection of parts in C1 stops B1; the identity of the parts is determined by the vision system when they arrive at the position C1. R1 gets part only if it can be assembled or temporary stored in ST<sub>i</sub>. Otherwise the part is left in B1. R2 gets completed assemblies from A1 or A2 and places them on B2 and B3 respectively.

2.3.2. System and Task Modelling

1) **System taxonomy.** The components of the assembly system are classed from a functional point of view: sensors, effectors, etc. This classification is useful to struc-

ture the factual knowledge of the assembly system: task state, component capabilities and relationships, etc. **Figure 2** shows the hierarchy concerning the assembly system of the example; the items in the lowest level of the hierarchy can be object instances of the upper concept (class).

2) **Workspace modelling.** In the example the following sites are defined: CONV1, CONV2 and CONV3 are the sites associated to the place where the parts stops in front of the robot; GRIP1 and GRIP2 are associated to the grippers of R1 and R2 respectively. The storing table has four sites: ST<sub>i</sub> (i = 1, ..., 4); they can be managed by the macro site ST. The sites associated to assembly tables are ASSB1 and ASSB2.

3) **Flow of parts.** The FPG shown in **Figure 5** describes the flow of parts required in the assembly task; the sites defined in the model are related by the operations whose outcome is, mainly, the transferring the parts between the sites. Operations may only modify the properties of the part into a site; as an example notice that the operation Ident-Loc does not transfer the part to another site but it changes the attributes of the unknown part.

Operations may also put parts into the flow model or drawn out parts (or products) from the model. The FPG for this example is defined as follows:

$$\begin{aligned}
 F &= (G, \text{SITES}, \text{OPER}, \lambda, \phi) \text{ where } G = (V, A) \text{ with} \\
 V &= \{v_0, v_1, v_2, v_3, v_4, v_5, v_6, v_7, v_8, v_9, v_{10}, v_{11}\} \\
 A &= \{(v_0, v_1), (v_1, v_2), (v_1, v_1), (v_1, v_3), (v_3, v_4), (v_4, v_3), \\
 &(v_3, v_5), (v_3, v_6), (v_5, v_7), (v_6, v_7), (v_7, v_8), (v_7, v_9), (v_8, v_{10}), \\
 &(v_8, v_{11})\} \\
 \text{SITES} &= \{\text{conv1}, \text{conv2}, \text{conv3}, \text{grip1}, \text{grip2}, \text{assb1}, \\
 &\text{assb2}, \text{st1}\} \\
 \varepsilon &= \{\text{input1}, \text{output1}, \text{output2}, \text{output3}\} \\
 \text{OPER} &= \{\text{B1.feed}(), \text{B1.advance}(), \text{R1.pick}(\text{conv1}), \text{R1.} \\
 &\text{store}(\text{st}), \text{R1.recover}(\text{st}), \text{R1.assembly}(\text{assb1}), \text{R1.assembly}(\text{assb2}), \\
 &\text{R2.pick}(\text{assb1}), \text{R2.pick}(\text{assb2}), \text{R2.place-}
 \end{aligned}$$

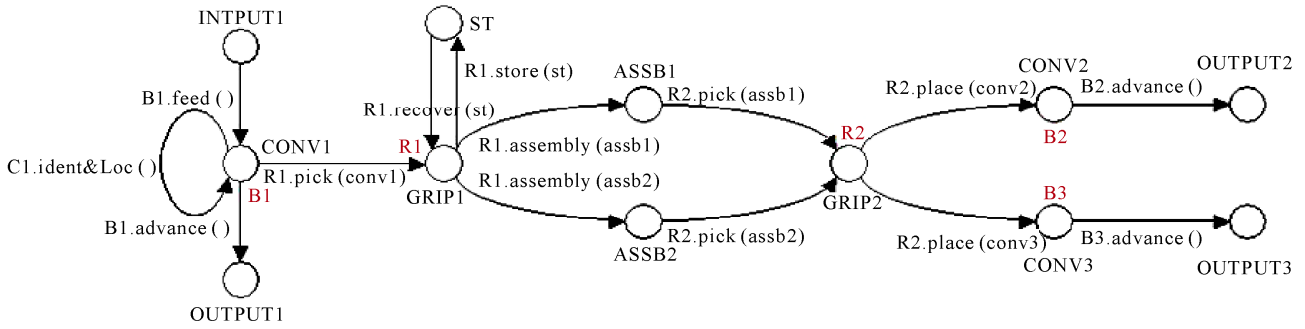


Figure 5. Flow of parts graphs for the assembly cell.

(conv2), R2.place2(conv3), B2.advance(), B3.advance(), C1.ident&Loc() }

$\lambda = \{(v_0, \text{input1}), (v_1, \text{conv1}), (v_2, \text{output1}), (v_3, \text{grip1}), (v_4, \text{st1}), (v_5, \text{assb1}), (v_6, \text{assb2}), (v_7, \text{grip2}), (v_8, \text{conv2}), (v_9, \text{conv3}), (v_{10}, \text{output2}), (v_{11}, \text{output3})\}$

$\phi = \{(v_0, v_1), \text{B1.feed}(), ((v_1, v_2), \text{B1.advance}()), ((v_1, v_3), \text{R1.pick}(\text{conv1})), ((v_3, v_4), \text{R1.store}(\text{st})), ((v_4, v_3), \text{R1.retrieve}(\text{st})), ((v_3, v_5), \text{R1.assembly}(\text{assb1})), ((v_3, v_6), \text{R1.assembly}(\text{assb2})), ((v_5, v_7), \text{R2.pick}(\text{assb1})), ((v_6, v_7), \text{R2.pick}(\text{assb2})), ((v_7, v_8), \text{R2.place}(\text{conv2})), ((v_7, v_9), \text{R2.place}(\text{conv3})), ((v_8, v_{10}), \text{B2.advance}()), ((v_8, v_{11}), \text{B3.advance}()), ((v_1, v_1), \text{C1.ident\&Loc}())\}$

### 3. Distributed Software Design

This section deals with the design of the distributed software that implements the task controller of a manufacturing system. First the modularisation of the task model is presented, and then the resulting partition is taken for implementing the agents [14].

#### 3.1. Task Model Partition

The task model must be decomposed into subtasks in such manner that every subtask may be assigned to an agent; this decomposition is achieved by a partition of the FPG. Several strategies for obtaining sub-graphs from the FPG may be adopted according to the number of processors, the geographical distribution on the components, or the similarity of the sub-graphs.

In this work the strategy held for decomposing the graph is creating the maximal number of sub-graphs; each sub-graph must involve one active site. This approach allows defining agents capable to control one effector. A three-step algorithm is described below.

**Algorithm.** Partitioning the task model.

1) *Identify active and passive sites.* Let  $\text{Act} \subseteq \text{SITES}$ , the set of active sites and  $\text{Pass} \subseteq \text{SITES}$  the set of passive sites.

2) *Sub-model creation.* For every  $s \in \text{Act}$ , create a sub-graph  $g_k$  including  $s$  and its predecessors and successors. In  $g_k$  it is included an active site and several passive sites.

$\text{SG} = \{g_1, g_2, \dots, g_r\}$  is the set of sub-models.

3) *Simplification of sub-models.* The operations associated to the arcs of a graph  $g_k$  must be executed by the effectors or sensors associated to the active site of  $g_k$ . Thus the arcs labelled with other operations must be withdrawn from  $g_k$ ; consequently isolated vertex must be eliminated too.

The sites belonging to two or more  $g_k$  are called *interface sites*; they are in the boundary of the graph and they must be carefully managed because they are considered as shared resources.

**Example.** Consider the FPG of Figure 5; defining  $\text{Act} = \{\text{conv1}, \text{conv2}, \text{conv3}, \text{grip1}, \text{grip2}\}$ , and  $\text{Pas} = \{\text{assb1}, \text{assb2}, \text{st}\}$ , the decomposition procedure yields the sub-graphs depicted in Figure 6; so  $\text{SG} = \{g_1, g_2, g_3, g_4, g_5\}$ .

#### 3.2. Task Programming Framework

Once the task model is decomposed, each sub-model is used for defining an agent. Since every sub-model involves an actuator, it must be controlled by the corresponding agent avoiding situations in which two or more agents handle the same effector. The knowledge base of every agent corresponds to the set of rules associated to the arcs of the pertaining subtask.

The developed platform supports the programming of the agents that implement the subtasks of the system; this platform is based on JADE V1.2 (Java Agent Development framework) [15], which meets the standards of FIPA [16].

##### 3.2.1. Requirements

**Agent requirements.** Each agent

- has a unique name
- manages the corresponding subtask; it initialises and updates the contents of its sites and macro-sites.
- exchanges messages with other agents; it interprets the message and perform the requests such as provide information about the contents of a site or about the execution of an operation.
- coordinates with other agents for allocating shared resources.

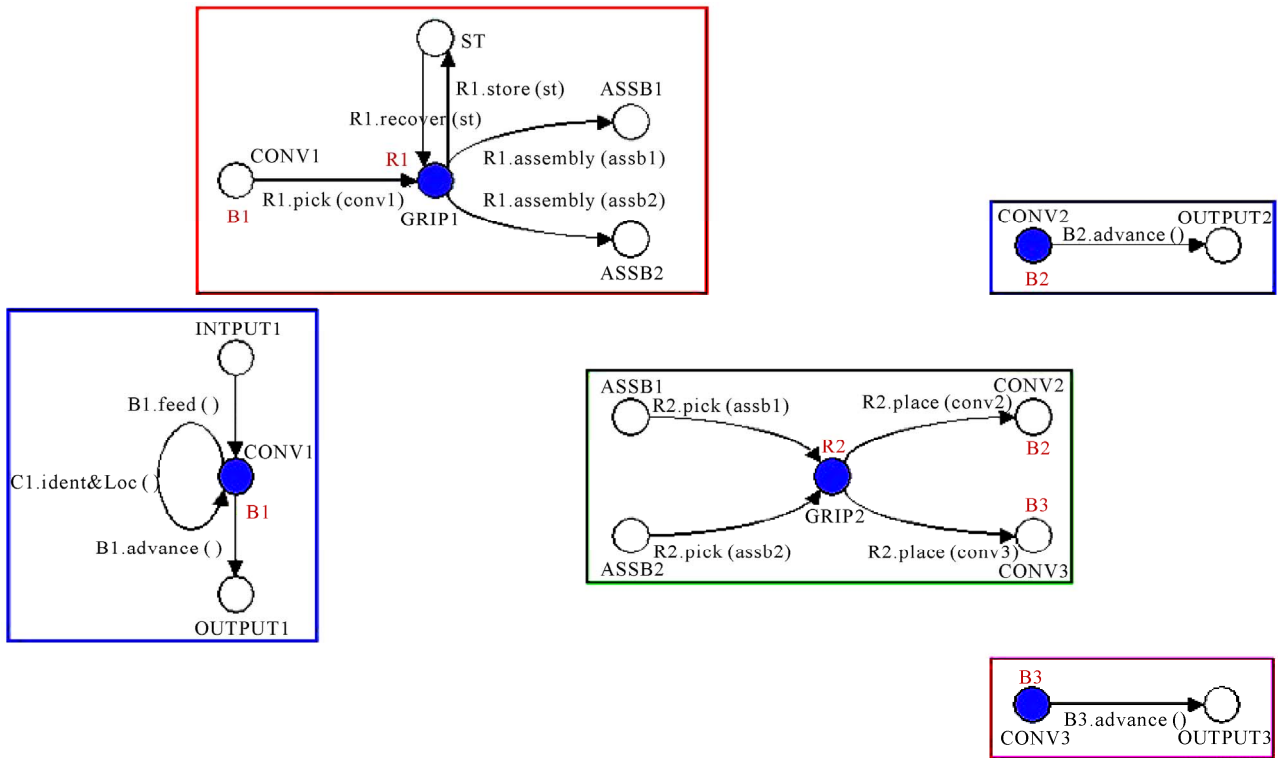


Figure 6. Graph decomposition.

- manages the associated devices (effectors and sensors involved in the subtask).

**Sites requirements.** Each site or macro-site

- has a unique name.
- provides facilities for managing its state, *i.e.* the initialising and updating of the associated attributes such as contents, position, sensor values, trace of the operations performed on it, etc.
- provides facilities for managing special features, such as flags for mutual exclusion, the agent names that share it, or the assembly patterns if it is a site for assembly.

**Devices requirements.** For every device (effector or sensor) in the system, one must create an interface module that allows consulting the device state and handling the messages representing actions requests or responses. Every module has a unique name.

**3.2.2. Definition of Classes**

**Component architecture.** The implemented components are integrated into a package organised in four sub-packages; it is shown in Figure 7. The sub-package MAgent.devices contains the classes *Site* and *Part*, MAgent.devices contains the class *Macrosite* and the classes related with the access of sites contained in the macrosite. Three kinds of sites are considered: *interface sites* (shared by two or more sub-tasks), *internal sites*, and *re-*

*mote sites* (non shared but belonging to other sub-tasks).

Figure 8 shows a detail of the above packages.

The sub-package MAgent.devices contains the class *Device* and other sub-packages. MAgent.devices.effectors and MAgent.devices.sensors contain sub-classes illustrated in Figure 9. This organisation is strongly suggested by the taxonomy of the manufacturing system obtained during the modelling stage.

**Agent class.** The class *AgentBase* is an abstract class that specialises the class *Agent* from JADE. This class em-

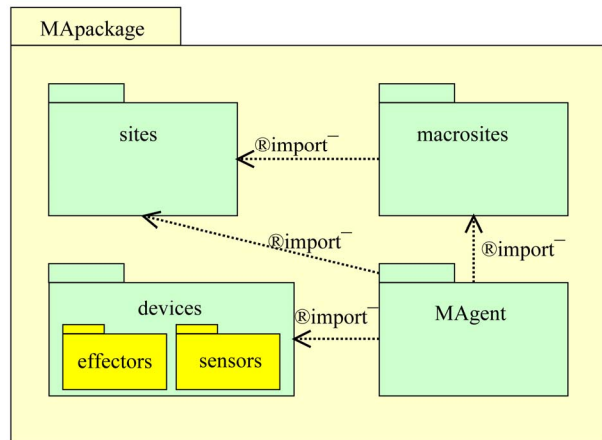


Figure 7. Component organisation.

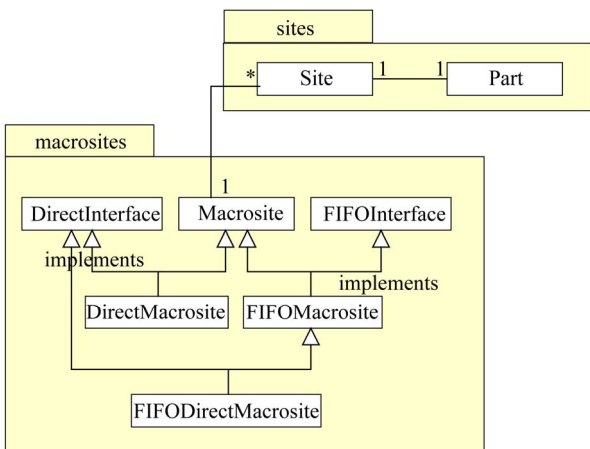


Figure 8. Relationships among sub-packages.

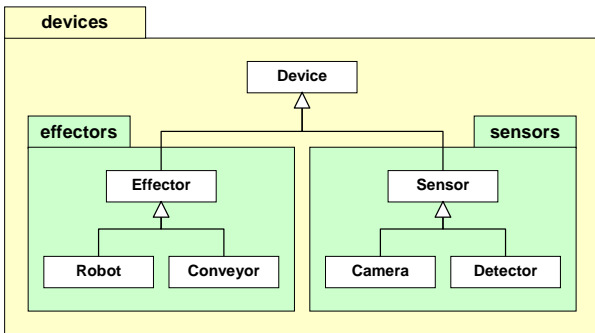


Figure 9. Devices sub-package.

beds the behaviour of a *generic agent* that manages all the activities related to a subtask. Every agent must be programmed by extending *AgentBase*, declaring the knowledge of the corresponding subtask, and instancing the extended subclass. The generic agent provides the facilities for the management of sites and the handling of messages related with the manufacturing task and messages for interacting with other agents; interaction among agents is performed through four kinds of messages regarding information requests and sending about sites, and request/confirmation on the use of sites shared by two or more agents.

**Agent programming.** The programming of each agent is done based in the information obtained from the corresponding subtask graph.

### 3.2.3. Component Identification

The first step for programming an agent is to identify the rules, the sites, the actions, and devices regarding the subtask. For every site in the graph one must declare which agent shares this site, and define if it is a remote site. If a site is used for assembly declare the assembly pattern to follow during the execution of the task. When an interface site is associated to sensor, only one agent must manage

| Operation              | Ident-Loc   |
|------------------------|---|
| <b>Description</b>     | If during the feeding a part is detected, the conveyor 1 stops and the part is identified. The sensor value is reset.                       |
| <b>Related sites</b>   | Conv1   |
| <b>Related devices</b> | Conveyor1, Vision system, Detector.   |
| <b>Pre-conditions</b>  | Conv1 is empty, the state of conveyor is FEEDING, the presence sensor is ON, and the subtask state is INITIAL.                              |
| <b>Actions</b>         | Stop Conveyor1. Identify and locate the part with the vision system. Reset sensor.  |
| <b>Post-conditions</b> | Conv1 holds a part, the state of Conveyor1 is STOP, the vision system is IDLE, and the sensor is OFF. The state of the task is HOLDINGPART. |

Figure 10. Frame for a rule description.

such sensor.

### 3.2.4. Building the Rule Base

The rules define the strategy of the controller. Every agent has a small set of rules corresponding to the arcs of the sub-graph. Before the writing of the rules it is convenient to enumerate all the information regarding each rule. This may be systematised by the filling of frames, such as shown in **Figure 10** for the operation Ident-Loc of the sub-task 1.

The Java coded rules for the agent coordinating the sub-task 3 is given below.

```
// implementing rules for the agent named
// ManufacturingAgent3
void rules()
{
    // declaring reference to site for place //the
    part
    Site destinationPart = null;
    // GRIP2.TAKE1: Grip2 takes the assembled //part
    from ASSB3
    if(assb3.getSite(0).isAssembled() &&
    grip2.isEmpty() &&
    conv2.isEmpty() && destinationPart==null)
    {
        grip2.setContent(assb3.getContent(0));
        assb3.removeContent(0);
        sendUpdateMessages(grip2);
        sendUpdateMessages(assb3);
        //
        destinationPart = conv2;
    }
    // GRIP2.TAKE2: Grip2 takes the assembled part
    // from ASSB4
    if(assb4.getSite(0).isAssembled())&&
    grip2.isEmpty() &&
    conv3.isEmpty() && destinationPart==null)
```

```

    {
        grip2.setContent(assb4.getContent(0));
        assb4.removeContent(0);
        sendUpdateMessages(grip2);
        sendUpdateMessages(assb4);
        destinationPart = conv3;
    }
// GRIP2.TAKE: Grip2 places the took part in
// CONVin takes the assembled part from ASSB4
    if(!grip2.isEmpty() && destination-
Part.isEmpty() && destinationPart!=null)
    {
        destination-
Part.setContent(grip2.getContent());
        grip2.removeContent();
        sendUpdateMessages(destinationPart);
        sendUpdateMessages(grip2);
        destinationPart = null;
    }
}

```

### 3.3. Implementation Issues

This method has been demonstrated through the software implementation of several case studies. The software has been written in Java using the JADE framework for supporting the agent definition and task interaction.

The distributed software has been tested on several personal computers interconnected through a local area network; every agent was assigned to single PC.

For executing the controller, first the JADE framework is initiated in a computer where an agent may be executed, and then the rest of the agents are started in their computers.

For every agent the interaction with a device has been simulated through a visual interface on screen; during a test, users simulate the devices response to commands sent by the controller through the keyboard.

### 4. Conclusions

In this work a method to develop distributed software for manufacturing control systems has been presented. One important issue of the proposed method is the modelling stage; the specifications are transformed into graphical models that contribute to reduce the problems due to ambiguities and incompleteness. The partitioned task model leads to obtain systematically the knowledge of every agent in the control software, allowing modifying easily the strategies of the subtasks; this feature is useful when task reprogramming is needed.

The programming methodology takes advantage of the JADE framework facilities, which permit to the defined agents to be executed into different kinds of platforms.

The case studies have been tested by simulating the manufacturing system devices through the console computers; however a real application could be developed by the implementation of the interface modules that communicate the controller with the actual devices.

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