

The Improvement of Relay Selection Schemes in Two-Hop Cellular Networks

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Abstract

In wireless networks, cooperative relaying could improve throughput by exploiting diversity. In order to reduce the amount of feedback for the channel gain, a semi-distributed scheme based on the relay feasible condition is proposed. Each relay node can measure its backward and forward channel gains. If both the channel gains are larger than a pre-defined threshold, this relay node is feasible. The final decision on the best relay selection is still given by the base station. Besides, the switch-and-examine relay selection scheme, which selects the first feasible relay node, is also investigated. Simulation results are presented to illustrate the advantage of two proposed schemes.

Keywords

Relay Selection, Cooperative Relaying, Amplify-and-Forward, Cellular Networks

1. Introduction

In next-generation wireless networks, cooperative relaying can be applied to extend the coverage and mitigate the signal fading arising from multi-path propagation or blocks [1] [2] [3] [4] [5]. The main idea is that single-antenna mobiles in a multi-user scenario can “share” their antennas in a manner that creates a virtual multiple-input multiple-output (MIMO) system.

In [6], the authors have proposed a simple cooperative diversity method, which selects the best relay from available relays and then uses the best relay for cooperation between the source and the destination. This distributed relay selection method requires no topology information and is based on local measurements of the instantaneous channel conditions. Moreover, this scheme can achieve the same diversity-multiplexing tradeoff as achieved by more complex protocols [7] [8] [9], where coordination and distributed space-time coding for

available relay nodes are required.

Generally, relay node selection in two-hop cellular networks [10] is much more manageable than that in ad hoc networks, mainly due to the presence of a central node (namely base station) with much more functionality and intelligence. Nevertheless, relay node selection is still a non-trivial issue in two-hop cellular networks since there will often be many candidate relay nodes for the source node and the failure to choose the optimal relay node would impair the overall performance improvement.

In [11], a centralized relay selection scheme is employed in the downlink of a cellular network. The base station uses the feedback to get the channel gain of every link. Then the base station calculates the available data rate through direct link (from base station to the mobile user) and triangular link (see in **Figure 1**) and determines the best link for each user. However, the channel gain feedback can cause a lot of burden to the system, and the base station applies the triangular link as long as the channel gain from the source node to the relay node is better than that of the direct link *i.e.* $\alpha_{si} > \alpha_{sd}$. The condition is not accurate.

2. System Model

We consider the downlink in a cellular network with cooperative relaying. A single broadband channel is shared by all users in TDMA manner. The system under consideration consists of a source node, K relay nodes, and N destination nodes. Throughout this paper, we consider half-duplex amplify-and-forward (AF) [3] relaying, where data transmission from the source to one destination requires two channel hops and two non-overlapping time slots. The transmission in the time domain is on a frame-by-frame basis. Each frame consists of two consecutive time slots. The link gain may integrate the effects from both propagation path loss and fading, and is varying independently among users, so that it is approximately unchanged during each frame interval.

For simplicity, we firstly consider this case, the source node “ s ” transmits information to a destination node “ d ” with the help of a relay node “ i ”. In the first time slot, the source node transmits to the destination as well as the relay node. In this phase, the signals received at the destination and the relay are respectively given by

$$y_{si} = \alpha_{si}x + n_{si} \quad (1)$$

$$y_{sd} = \alpha_{sd}x + n_{sd} \quad (2)$$

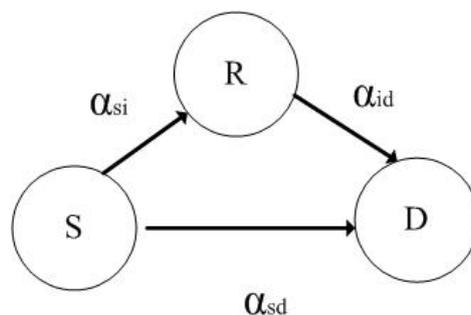


Figure 1. Cooperative relaying: triangular model.

where x, y_{sd}, y_{si} denote the transmitted signal, the signal received at the destination and the signal received at the i th relay node, respectively. α_{si} and α_{sd} are channel coefficients of the source-relay and source-destination channels. n_{si} and n_{sd} denote the background noise at the relay node i and destination node d , respectively, which are independent identically distributed (i.i.d) complex Gaussian random variables with a common variance σ_n^2 .

During the second time slot, while the source node is idle, the relay node amplifies and forwards its received signal from the source to the destination. The received signal at destination node d from relay node i is given by

$$y_{id} = \alpha_{id}\lambda_i y_{si} + n_{id} = \alpha_{id}\lambda_i\alpha_{si}x + \alpha_{id}\lambda_i n_{si} + n_{id} \tag{3}$$

where α_{id} is the channel gain between the relay node and the destination node. $n_{id} \sim CN(0, \sigma_n^2)$ denotes the noise at the destination. λ_i is an amplification factor of the i th relay, which is used to guarantee the transmission power of the relay node and satisfies

$$\lambda_i^2 = \frac{P_i}{|\alpha_{si}|^2 P_s + \sigma_n^2} \tag{4}$$

where P_s and P_i denote the transmission power of the source node and the relay node, respectively.

3. Capacity Analysis of the Cellular Network with Cooperative Relaying

In this section, we analyze the capacity of the cellular network, which consists of one source, K relay nodes, and N destination nodes. For simplicity, we first analyze the scenario, which consists of only a single destination node, and then extend the results to the case of multiple destination nodes.

If α_{si}, α_{id} and α_{sd} are known and relay i is chosen for relaying, the channel combining both the direct path and the relay path can be modeled as an equivalent one-input, two-output complex Gaussian noise channel, which has the maximum average mutual information given by [3]

$$I_{AF}^i = \frac{1}{2} \log \left(1 + \frac{P_s}{\sigma_n^2} \left(\beta_{sd} + \frac{\beta_{si}\beta_{id}P_i}{\beta_{si}P_s + \beta_{id}P_i + \sigma_n^2} \right) \right) \tag{5}$$

where $\beta_{si} = |\alpha_{si}|^2$, $\beta_{sd} = |\alpha_{sd}|^2$, and $\beta_{id} = |\alpha_{id}|^2$. To focus on the idea of relay selection, we assume equal power allocation between any pair of the source and relay nodes, *i.e.* $P_s = P_i = P_{\max} / 2 = P$, and define the transmit signal to noise ratio (SNR) as $\gamma = P / \sigma_n^2$. Thus, Equation (5) can be rewritten as

$$I_{AF}^i = \frac{1}{2} \log_2 \left(1 + \gamma \left(\beta_{sd} + \frac{\gamma\beta_{si}\beta_{id}}{\gamma(\beta_{si} + \beta_{id}) + 1} \right) \right) \tag{6}$$

Therefore, the maximum capacity would be attained when the relay with the largest A_i is selected as

$$A_i = \frac{\gamma\beta_{si}\beta_{id}}{\gamma(\beta_{si} + \beta_{id}) + 1} \tag{7}$$

which results in a capacity of

$$\max \{I_{AF}^i\} = \frac{1}{2} \log_2 \left(1 + \gamma(\beta_{sd} + \max_i \frac{\gamma\beta_{si}\beta_{id}}{\gamma(\beta_{si} + \beta_{id}) + 1}) \right) \quad (8)$$

If no relaying is applied and both time slots in one frame contribute to the source node’s transmissions, the achievable average mutual information for the direct transmission can be calculated as

$$I_D = 2 \times \frac{1}{2} \log_2(1 + \beta_{sd}\gamma) = \log_2(1 + \beta_{sd}\gamma) \quad (9)$$

Next, the conditions for selecting a feasible relay node are considered. A relay node is feasible if user relaying via it can provide better capacity performance than direct transmission, that is, $I_{AF}^i > I_D$. Define Γ_f as the set of feasible relay nodes. According to the Lemma 2 in [12], the sufficient condition that any relay node i belongs to the feasible set is

$$\beta_{si} > C_s \text{ and } \beta_{id} > C_s, \quad \forall i \in \Gamma_f \quad (10)$$

where

$$C_s = \left(1 + \sqrt{1 + \frac{1}{\gamma C_n}} \right) C_n, \quad C_n = \beta_{sd}(1 + \beta_{sd}\gamma) \quad (11)$$

We call the above condition “the relay feasible condition”. The condition for determining the feasibility of relay nodes ensures that user relaying via the node can achieve a larger channel capacity than direct transmission [12]. Moreover, the condition significantly reduces the search time for the best relay node. Note that we define the search time as the number of nodes searched during one execution of the search algorithm. As a result, the computational complexity can be effectively reduced. But in [11], the relay node is considered to be feasible as long as $\beta_{si} > \beta_{sd}$. Since $\beta_{sd} < C_s$, for the relay selection scheme in [11], the searching time will be increased.

Combining Equation (8) and Equation (9), the maximum average mutual information for a pair of cooperating users is equal to

$$C = \max \left\{ I_D, \max_{i \in \Gamma_f} (I_{AF}^i) \right\} \quad (12)$$

maximum average mutual information for N pairs of cooperating users. It can be denoted by

$$C_{multiple_D} = \sum_{k=1}^N \max \left\{ I_{D,k}, \max_{i \in \Gamma_f} \{I_{AF,k}^i\} \right\} \quad (13)$$

4. Two Proposed Relay Selection Schemes

4.1. Switch-and-Examine Relay Selection Scheme

Although the relay selection scheme in [6] can provide the multiplexing and diversity trade-off as distributed space-time coding, it may lead to the packet collision. In addition, all relay candidates must make works like the path estimations during every transmission which causes the power consumption of relay nodes.

To overcome the issues, a sub-optimal relay selection scheme is proposed in [13], where the relay node is selected if its link quality is above a certain threshold which is set to satisfy the required performance and the selected link is maintained unless its link quality does not fall below the threshold. Note that the relay nodes work at Decode-and-Forward (DF) mode in [13].

Nevertheless, in the scheme in [13], there is no exact definition about the target threshold. Moreover, if the first $K - 1$ relay nodes are not acceptable, the last K th node will be chosen as the relay node without comparing to the target threshold. But when the backward and forward channels of the K th node suffer deep fading, the achievable capacity with the help of the relay node may be lower than that with direct transmission. Besides, the scheme in [13] is proposed for the case of DF mode at the relay node. Thus, we propose a novel relay selection scheme for the half-duplex AF relaying.

We assume there are K relay nodes in the system. In addition, it is assumed that an arbitrary selected or tested relay node can know its channel gains both source-relay and relay-destination link using an additional feedback link or a clear-to-send (CTS) packet from the destination [6]. Firstly, the channel gain between the source and the first relay node is estimated, where we denote it as $\beta_{s,l}$. C_s is the target threshold as defined in Equation (11). If the current path is acceptable (*i.e.* $\beta_{s,l} > C_s$) at the relay, the tested relay node will request the channel state information (CSI) to the destination node. When $\beta_{l,d} > C_s$, we decide it as an acceptable path and thus have the final channel gains of backward and forward channel as $\beta_1 = \beta_{s,l}$ and $\beta_2 = \beta_{l,d}$, respectively. If none of the relays satisfies above condition, the direct link will be used. We show the mode of operation in **Figure 2**.

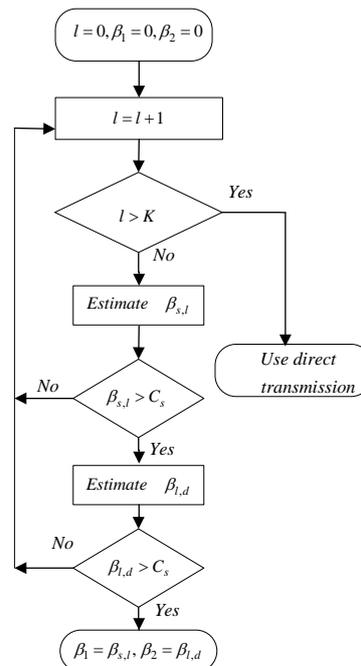


Figure 2. The block diagram of the modified switch-and-examine relay selection scheme.

As can be seen from the above scheme, setting the target threshold as C_s can ensure the achievable capacity via the selected relay is larger than that with direct transmission. If there is no relay to meet the requirement about the threshold, direct transmission link will be used. Therefore, the proposed scheme outperforms the scheme in [13].

4.2. The Semi-Distributed Relay Selection Scheme

As analyzed above, switch-and-examine relay selection scheme is a distributed scheme in which each relay can make decision on its feasibility individually. However, this scheme can't maximize the system capacity because the selected relay is not optimal and it is the first one which satisfies the condition (*i.e.* $\beta_{s,i} > C_s$ and $\beta_{t,d} > C_s$).

In order to deal with above problems, we propose a semi-distributed relay selection scheme for the downlink of cellular networks. The proposed relay selection scheme can be described as follows.

Step 1. The decision method of feasible set is the same as the idea of [12]. Before a source node initiates packet transmission to the destination, some hand-shaking signals have to be exchanged, such as request to send (RTS) and CTS in ad hoc networks. Similarly, some pilot signals need to be exchanged in cellular networks before data transmission. We borrow the names RTS and CTS to represent the hand-shaking signals before data transmission. After receiving the RTS, the destination node estimates the channel gain, β_{sd} , from the source node and calculates the threshold C_s (as defined in Equation (11)) based on the estimation. The destination feeds back a CTS as an acknowledgement, which includes the information on the threshold C_s .

Step 2. The relay nodes, who received both RTS and CTS, can estimate the channel gain, β_{si} and β_{id} . Then the relay nodes compare the channel gains with the threshold C_s to determine its feasibility. If the relay node is feasible, then it will calculate its A_i which is defined as Equation (7) and notify the source node of its ID and its A_i .

Step 3. After receiving all notifications from the feasible relay nodes, the source node generates a feasible set for each destination node. The source node sorts the destination nodes in ascending order based on the estimates of β_{sd} and the destination node with bad channel conditions will achieve higher priority to choose the relay node.

Step 4. For a given destination node, the relay node which has the maximal A_i will be chosen in its corresponding feasible set. Note that in [12], the relay node is randomly chosen from its feasible set.

Step 5. The source node removes the selected destination node's feasible set and removes the selected relay node from all feasible sets.

Step 6. The source node repeats the previous steps until there is no non-zero feasible set. The source node broadcasts the relay node allocation to all nodes. Then it will start data transmission.

Based on the above procedure of relay selection, we see that the proposed

semi-distributed scheme can achieve higher capacity than that of the relay selection scheme in [12]. This is due to the fact that for the scheme in [12], one relay node is randomly chosen from each destination's feasible set, but for the proposed scheme, the optimal relay node which can help the source to achieve the maximal capacity is selected from the feasible set. Moreover, in this relay selection scheme, each relay node can determine its feasibility according to the threshold C_s individually and there is no need for the source node to know every channel gain by information feedback, which can significantly reduce the system overhead. The final relay allocation is made by the source node in centralized manner. Therefore, compared to the scheme in [11], the proposed scheme has stringent constraints on system overhead.

5. Simulation Results

In this section, we first compare the switch-and-examine relay selection scheme with the scheme in [13]. Consider a circle centered at the origin of the $x-y$ plane with radius of 50 meters. The source node is located at the center of the circle, while all the other nodes are uniformly distributed in the circle. The channel between two nodes is $\alpha_{ij} \sim CN(0, 1/d^v)$, where d is the distance between two nodes and v is the path loss exponent, which equals 2.5 in our simulation. The channel gain integrates the effects from both propagation path loss and fading, and varies slowly in time. The background noise is i.i.d complex Gaussian random variable. The transmission power is represented through SNR at the transmitter end.

Firstly, we consider a simple case composed of one source node, 99 relay nodes and one destination node. The destination is located at the edge of the coverage area and the relay nodes are uniformly distributed in the circle. **Figure 3** shows the performance comparison of the switch-and-examine relay selection

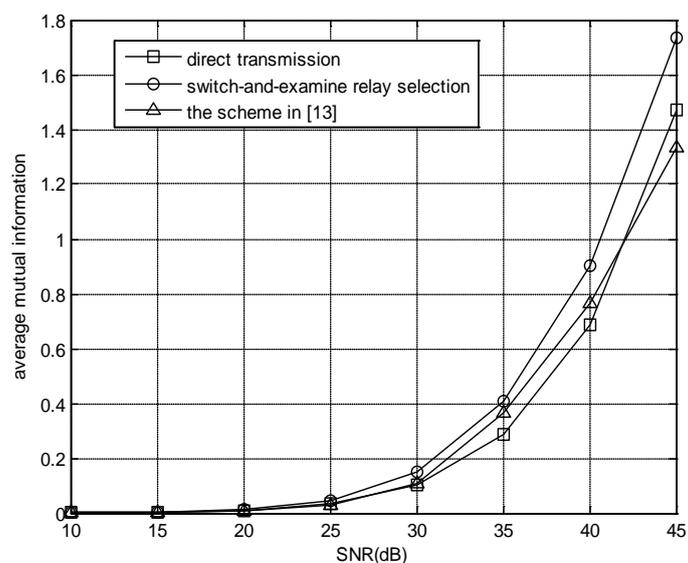


Figure 3. Performance comparison of switch-and-examine relay selection scheme and the scheme in [13].

scheme and the scheme in [13]. With the increase of transmit SNR, switch-and-examine relay selection scheme always gains higher average mutual information than that with the scheme in [13].

Figure 4 illustrates average mutual information versus the transmit SNR of five relay selection schemes. It can be shown that the proposed semi-distributed scheme gains higher capacity than the relay selection scheme in [12] and switch-and-examine scheme. This verifies the analysis in Section IV. Besides, it can be found that the semi-distributed scheme gets almost equivalent capacity with that of the scheme in [11]. **Figure 4** also shows that the direct transmission has relatively poor performance than other schemes. This is because direct transmission does not benefit from cooperative diversity at high SNR.

Figure 5 shows the proposed scheme exhibits better average search time for the best relay than that of the scheme in [11]. For instance, at SNR = 40 dB, the average search time using the proposed scheme is nearly one sixth of those using the scheme in [11]. Combining **Figure 4** with **Figure 5**, we can observe that the semi-distributed relay selection scheme provides the same capacity gains with the scheme in [11], but need much lower computational complexity than the scheme in [11].

Figure 6 shows the case with multiple destination nodes. The source is located at the center and other 100 nodes are uniformly distributed in the circle. One half of the 100 nodes are chosen as the destination nodes and the other half are considered as the relay nodes. It can be found that the semi-distributed scheme outperforms other schemes.

6. Conclusion

In this paper, we propose a semi-distributed relay selection scheme on the relay

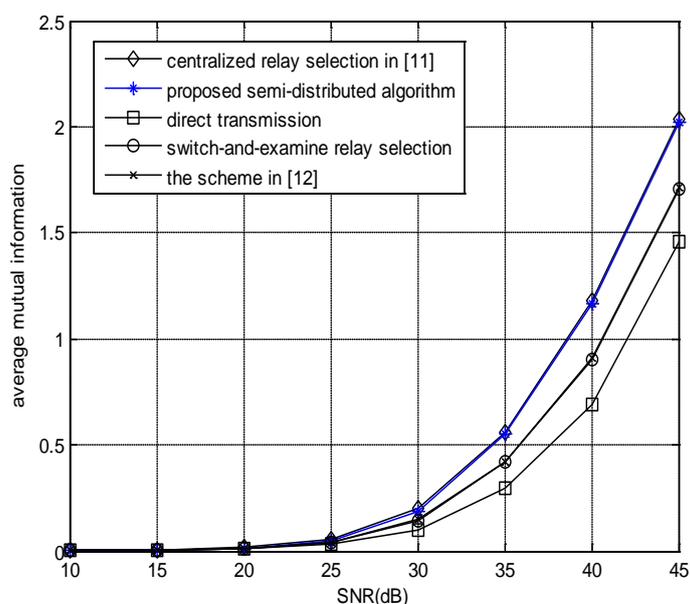


Figure 4. Average mutual information of five relay selection schemes with different SNR (one destination).

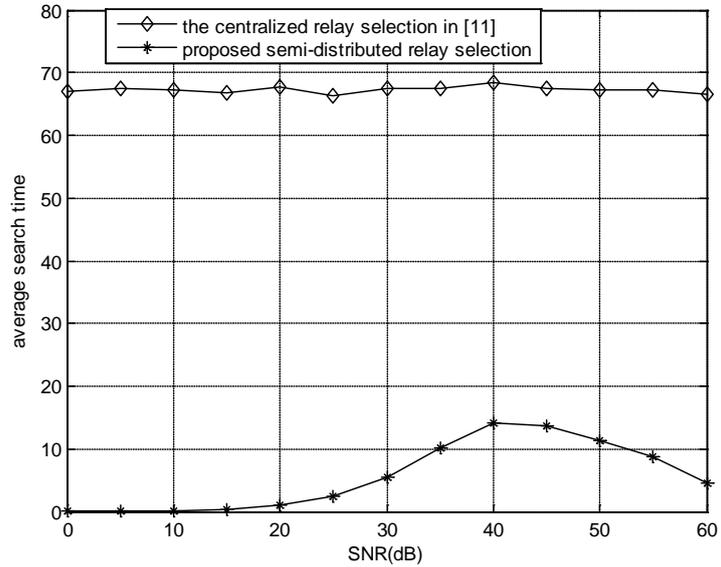


Figure 5. Computational complexity of two relay selection schemes with different SNR.

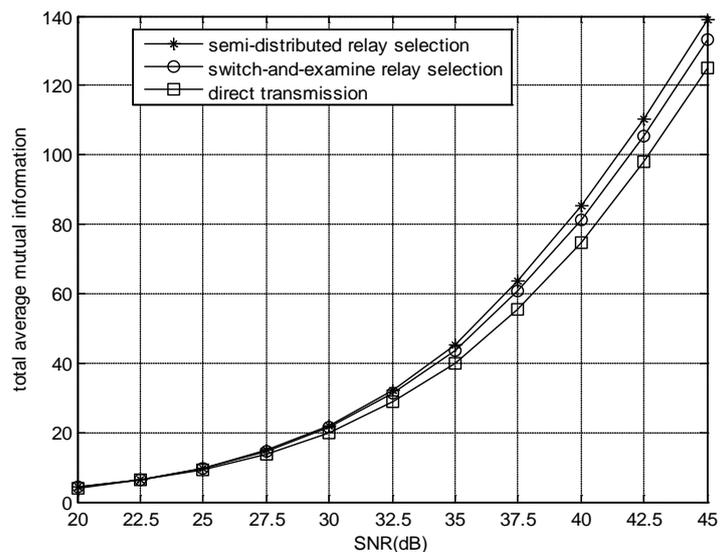


Figure 6. Total average mutual information of different relay selection scheme.

feasible condition for the downlink of cellular networks. Simulation results show that the semi-distributed scheme can achieve the nearly equivalent capacity compared with the scheme in [11], but the former has significantly reduced computational complexity. We also discuss the switch-and-examine relay selection scheme, which can reduce the complexity and power consumption at the relays. Simulation results validate our analysis and show that the semi-distributed scheme outperforms the switch-and-examine scheme. Besides, the two relay selection schemes can achieve much more capacity than direct transmission at high SNR because direct transmission does not benefit from cooperative diversity.

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