

Solution of Matrix Game with Triangular Intuitionistic Fuzzy Pay-Off Using Score Function

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ABSTRACT

Using score function in a matrix game is very rare. In the proposed paper we have considered a matrix game with pay-off as triangular intuitionistic fuzzy number and a new ranking order has been proposed using value judgement index, available definitions and operations. A new concept of score function has been developed to defuzzify the pay-off matrix and solution of the matrix game has been obtained. A numerical example has been given in support of the proposed method.

Keywords: Triangular Intuitionistic Fuzzy Number; Matrix Game; Value Judgement Index; Score Function

1. Introduction

Game theory is the way to handle the problems where two conflicting interests situation exist. But in modern society a lot of problems exist which cannot be explained in simple crisp sense e.g. there may be the situation where pay-offs are not known precisely. In such cases fuzzy mathematics is a tool to handle such situation. Fuzziness in matrix games can appear in many ways but two classes of fuzziness seem to be very natural. These two classes of fuzzy matrix games are referred to as matrix games with fuzzy goal [1] and matrix games with fuzzy pay off [2]. In recent times much attention has been drawn to interval valued game, Nayak and Pal [3-5], Narayanan [6], Nishizaki [1]. In practical situations the pay-offs are given with in certain ranges rather than as an exact number. These uncertain situations are overcome when we use interval numbers as pay-offs. An interval number is an extension of a real number and also a subset of a real line \mathcal{R} , Moore [7]. Zimmermann [8] shows that α cut of a fuzzy number is an interval number. The method of solution of a matrix game using interval numbers was already established, Nayak and Pal [5]. In Narayanan [6], probability and possibility approaches have been used to solve a 2×2 interval game but no certain distribution function has been used. Moreover, reduction of an $m \times n$ game to a 2×2 sub game is a basic problem in an interval game. In the dominance method [3], if the convex combination of any two rows(columns)of a pay-off matrix is dominated by the third row (column), it indicates that the third move of the

row (column) of the player will be an optimal move but we are not certain as to which one of the first two moves will be an optimal one. This disadvantage is overcome through the graphical method [4]. But it may be the situation where the players can estimate the approximate pay-off values with some degree but there exist a hesitation. Such situations are handled by intuitionistic fuzzy (IF) numbers. Atanassov [9] first introduced the concept of IF-set where he explained an element of an IF-set in respect of degree of belongingness, degree of non-belongingness and degree of hesitancy. This degree of hesitancy is nothing but the uncertainty in taking a decision by a decision maker (DM). Atanassov [10] first described a game using the IF-set. Li and Nan [11] considered the matrix games with pay-offs as IF-sets. Seikh, Nayak and Pal [12] considered a bi-matrix game where they used IF-set. In this paper we have considered a matrix game where the pay-off elements are considered as triangular intuitionistic fuzzy number (TIFN). Nan, Li and Zhang [13] considered such TIFN as pay-off elements of the matrix and described the arithmetic operation and cut sets. In this paper, we have made a ranking order of the TIFN based on value judgement index and deviation index of membership and non-membership functions. A score function approach has been described to defuzzify the matrix. The numerical problem is a real life voting share problem and establishes the theory on strong ground.

The paper is organized as follows: In Section 2, basic definitions of intuitionistic fuzzy set is given intuitionistic fuzzy number, TIFN and score function are defined and arithmetic operations are described. In Section 3,

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matrix game with TIFN pay-off, pure and mixed strategy have been described. In Section 4, numerical example is given. In Section 5 conclusion has been drawn.

2. Intuitionistic Fuzzy Sets

Here we are to introduce first some relevant basic preliminaries, notations and definitions of IFS, in particular the works of Atanassov [9,14].

Definition 1: Let $X = \{x_1, x_2, \dots, x_n\}$ be a finite universal set. An Atanassov's intuitionistic fuzzy set (IFS) in a given universal set X is an expression A given by

$$A = \{ \langle x_i, \mu_A(x_i), \nu_A(x_i) \rangle : x_i \in X \} \tag{1}$$

where the functions

$$\mu_A : X \rightarrow [0,1]; \quad x_i \in X \rightarrow \mu_A(x_i) \in [0,1]$$

and $\nu_A : X \rightarrow [0,1]; \quad x_i \in X \rightarrow \nu_A(x_i) \in [0,1]$ define the degree of membership and the degree of non-membership of an element $x_i \in X$ to the set $A \subseteq X$, respectively, such that they satisfy the following condition: for every $x_i \in X$

$$0 \leq \mu_A(x) + \nu_A(x) \leq 1.$$

Let

$$\pi_A(x_i) = 1 - \mu_A(x) - \nu_A(x)$$

which is called the Atanassov's [14] intuitionistic index of an element x_i in the set A . It is the degree of indeterminacy membership of the element x_i to the set A . Obviously, $0 \leq \pi_A(x_i) \leq 1$. If an Atanassov's IFS C in X has only an element, then C is written as follows

$$C = \{ \langle x_k, \mu_C(x_k), \nu_C(x_k) \rangle \}$$

which is usually denoted by $C = \{ \langle \mu_C(x_k), \nu_C(x_k) \rangle \}$ for short.

Definition 2: Let A and B be two Atanassov's IFSs in the set X . $A \subset B$ iff

$$\mu_A(x_i) \leq \mu_B(x_i) \text{ and } \nu_A(x_i) \geq \nu_B(x_i); \text{ for any } x_i \in X.$$

Definition 3: Let A and B be two Atanassov's IFSs in the set X . $A = B$ iff

$$\mu_A(x_i) = \mu_B(x_i) \text{ and } \nu_A(x_i) = \nu_B(x_i); \text{ for any } x_i \in X$$

Namely, $A = B$ iff $A \subset B$ and $B \subset A$.

Definition 4: Let A and B be two Atanassov's IFSs in the set X . The intersection of A and B is defined as follows:

$$A \cap B = \{ \langle x_i, \min(\mu_A(x_i), \mu_B(x_i)), \max(\nu_A(x_i), \nu_B(x_i)) \rangle | x_i \in X \}.$$

Definition 5: (Intuitionistic Fuzzy Number [15]): An intuitionistic fuzzy number (**Figure 1**) \tilde{A}^i is

- 1) an intuitionistic fuzzy subset of the real line;
- 2) normal, *i.e.* there exists $x_0 \in \mathfrak{R}$ such that

$$\mu_{\tilde{A}^i}(x_0) = 1 \text{ (so } \nu_{\tilde{A}^i}(x_0) = 0 \text{)};$$

- 3) convex for the membership function $\mu_{\tilde{A}^i}$ *i.e.*

$$\mu_{\tilde{A}^i}(\lambda x_1 + (1 - \lambda)x_2) \geq \min \{ \mu_{\tilde{A}^i}(x_1), \mu_{\tilde{A}^i}(x_2) \}; \quad \forall x_1, x_2 \in \mathfrak{R}, \lambda \in [0,1]$$

- 4) concave for the non-membership function $\nu_{\tilde{A}^i}$ *i.e.*

$$\nu_{\tilde{A}^i}(\lambda x_1 + (1 - \lambda)x_2) \leq \max \{ \nu_{\tilde{A}^i}(x_1), \nu_{\tilde{A}^i}(x_2) \}; \quad \forall x_1, x_2 \in \mathfrak{R}, \lambda \in [0,1].$$

In our discussion we consider an intuitionistic fuzzy number \tilde{A}^i as $\langle \mu_{ij}, \nu_{ij} \rangle$ where $\mu_{ij} = \mu_{\tilde{A}^i}(x_{ij})$ and $\nu_{ij} = \nu_{\tilde{A}^i}(x_{ij})$ as we consider it as ij th element of cost matrix.

2.1. Triangular Intuitionistic Fuzzy Number

The definitions and operations of TIFN given by Nan, Li and Zhang [13] are stated as follows:

Definition 6: An TIFN $\tilde{t} = \langle \langle \underline{t}, \bar{t} \rangle, \chi_{\tilde{t}}, \eta_{\tilde{t}} \rangle$ defined on the real number set \mathfrak{R} is an intuitionistic fuzzy set, whose membership and non-membership function are given by

$$\mu_{\tilde{t}}(x) = \begin{cases} \chi_{\tilde{t}}(x - \underline{t}) / (t - \underline{t}), & \text{if } \underline{t} \leq x < t \\ \chi_{\tilde{t}}, & \text{if } x = t \\ \chi_{\tilde{t}}(\bar{t} - x) / (\bar{t} - t), & \text{if } t < x \leq \bar{t} \\ 0, & \text{if } x < \underline{t} \text{ or } x > \bar{t}. \end{cases}$$

and

$$\nu_{\tilde{t}}(x) = \begin{cases} \{a - x + \eta_{\tilde{t}}(x - \underline{t})\} / (t - \underline{t}), & \text{if } \underline{t} \leq x < t \\ \eta_{\tilde{t}}, & \text{if } x = t \\ \{x - a + \eta_{\tilde{t}}(\bar{t} - x)\} / (\bar{t} - t), & \text{if } t < x \leq \bar{t} \\ 1, & \text{if } x < \underline{t} \text{ or } x > \bar{t}. \end{cases}$$

respectively, where the values $\chi_{\tilde{t}}$ and $\eta_{\tilde{t}}$ represent the maximum degree of membership and the minimum degree of non-membership, respectively, such that they satisfy the following condition: $0 \leq \chi_{\tilde{t}} \leq 1, 0 \leq \eta_{\tilde{t}} \leq 1$ and $0 \leq \chi_{\tilde{t}} + \eta_{\tilde{t}} \leq 1$.

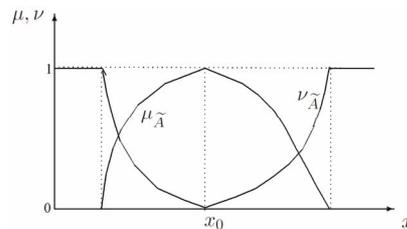


Figure 1. Membership and non-membership functions of \tilde{A}^i .

The hesitancy degree or the degree of indeterminacy membership of the element x to the TIFN \tilde{t} can be given as $\pi_{\tilde{t}}(x) = 1 - \mu_{\tilde{t}}(x) - \nu_{\tilde{t}}(x)$. Here χ_i and η_i represent respectively the confidence and non-confidence levels of the TIFN \tilde{t} .

Definition 7: Let us consider two TIFNs as $\tilde{t} = \langle (\underline{t}, t, \bar{t}); \chi_i, \eta_i \rangle$, $\tilde{s} = \langle (\underline{s}, s, \bar{s}); \chi_s, \eta_s \rangle$. With $\chi_i \neq \chi_s$, the arithmetic operations are defined as follows:

$$1) \tilde{t} + \tilde{s} = \langle (\underline{t} + \underline{s}, t + s, \bar{t} + \bar{s}); \chi_i \wedge \chi_s, \eta_i \vee \eta_s \rangle,$$

where “ \wedge ” and “ \vee ” represent min and max operators respectively.

$$2) \tilde{t} - \tilde{s} = \langle (\underline{t} - \underline{s}, t - s, \bar{t} - \bar{s}); \chi_i \wedge \chi_s, \eta_i \vee \eta_s \rangle$$

3)

$$\tilde{t} \times \tilde{s} = \begin{cases} \langle (\underline{t}\underline{s}, ts, \bar{t}\bar{s}) \chi_i \wedge \chi_s, \eta_i \vee \eta_s \rangle, & \text{if } \tilde{t} > 0 \text{ and } \tilde{s} > 0 \\ \langle (\underline{t}\bar{s}, ts, \bar{t}\underline{s}) \chi_i \wedge \chi_s, \eta_i \vee \eta_s \rangle, & \text{if } \tilde{t} < 0 \text{ and } \tilde{s} > 0 \\ \langle (\bar{t}\underline{s}, ts, \underline{t}\underline{s}) \chi_i \wedge \chi_s, \eta_i \vee \eta_s \rangle, & \text{if } \tilde{t} < 0 \text{ and } \tilde{s} < 0 \end{cases}$$

$$4) \tilde{t}/\tilde{s} = \begin{cases} \langle (\underline{t}/\underline{s}, t/s, \bar{t}/\bar{s}) \chi_i \wedge \chi_s, \eta_i \vee \eta_s \rangle, & \text{if } \tilde{t} > 0 \text{ and } \tilde{s} > 0 \\ \langle (\bar{t}/\underline{s}, t/s, \underline{t}/\underline{s}) \chi_i \wedge \chi_s, \eta_i \vee \eta_s \rangle, & \text{if } \tilde{t} < 0 \text{ and } \tilde{s} > 0 \\ \langle (\bar{t}/\underline{s}, t/s, \underline{t}/\bar{s}) \chi_i \wedge \chi_s, \eta_i \vee \eta_s \rangle, & \text{if } \tilde{t} < 0 \text{ and } \tilde{s} < 0 \end{cases}$$

$$5) \lambda \tilde{t} = \begin{cases} \langle (\lambda \underline{t}, \lambda t, \lambda \bar{t}) \chi_i, \eta_i \rangle, & \text{if } \lambda > 0 \\ \langle (\lambda \bar{t}, \lambda t, \lambda \underline{t}) \chi_i, \eta_i \rangle, & \text{if } \lambda < 0 \end{cases}$$

where λ is any real number.

$$6) \tilde{t}^{-1} = \langle (1/\bar{t}, 1/t, 1/\underline{t}) \chi_i, \eta_i \rangle$$

2.2. Cut Sets of TIFN

Definition 8: A (α, β) - cut set of $\tilde{t} = \langle (\underline{t}, t, \bar{t}) \chi_i, \eta_i \rangle$ is defined over a crisp subset of \mathfrak{R} and it is given as

$$\tilde{t}_{\alpha, \beta} = \{x : \mu_{\tilde{t}}(x) \geq \alpha, \nu_{\tilde{t}}(x) \leq \beta\},$$

where $0 \leq \alpha \leq \chi_i, \eta_i \leq \beta \leq 1$ and $0 \leq \alpha + \beta \leq 1$.

Definition 9: A α -cut set of $\tilde{t} = \langle (\underline{t}, t, \bar{t}) \chi_i, \eta_i \rangle$ is defined over a crisp subset of \mathfrak{R} and it is given as $\tilde{t}_{\alpha} = \{x : \mu_{\tilde{t}}(x) \geq \alpha\}$. Corresponding closed interval is given as $\tilde{t}_{\alpha} = [\underline{t} + \alpha(t - \underline{t})/\chi_i, t - \alpha(\bar{t} - t)/\chi_i]$.

Definition 10: A β -cut set of $\tilde{t} = \langle (\underline{t}, t, \bar{t}) \chi_i, \eta_i \rangle$ is defined over a crisp subset of \mathfrak{R} and it is given as $\tilde{t}_{\beta} = \{x : \mu_{\tilde{t}}(x) \leq \beta\}$. Corresponding closed interval is given as

$$\tilde{t}_{\beta} = \left[\frac{\{(1 - \beta)t + (\beta - \eta_i)\underline{t}\}}{(1 - \eta_i)}, \frac{\{(1 - \beta)t + (\beta - \eta_i)\bar{t}\}}{(1 - \eta_i)} \right].$$

Definition 11: Let $m(\tilde{t}_{\alpha})$ and $m(\tilde{t}_{\beta})$ be the mean

values of the intervals \tilde{t}_{α} and \tilde{t}_{β} respectively i.e.

$$m(\tilde{t}_{\alpha}) = [2t\alpha + (\chi_i - \alpha)(\underline{t} + \bar{t})]/2\chi_i$$

and

$$m(\tilde{t}_{\beta}) = [2t(1 - \beta) + (\beta - \eta_i)(\underline{t} + \bar{t})]/2(1 - \eta_i).$$

Then average index of the membership function $\mu_{\tilde{t}}(x)$ and the average index of the non-membership function $\nu_{\tilde{t}}(x)$ for the TIFN \tilde{t} are defined as

$$A_{\mu}(\tilde{t}) = \int_0^{\chi_i} m(\tilde{t}_{\alpha}) d\alpha = \chi_i(2t + \underline{t} + \bar{t})/4$$

$$A_{\nu}(\tilde{t}) = \int_{\eta_i}^1 m(\tilde{t}_{\beta}) d\beta = (1 - \eta_i)(2t + \underline{t} + \bar{t})/4$$

respectively. Now we will introduce deviation index of the membership function $\mu_{\tilde{t}}(x)$ and non-membership function $\nu_{\tilde{t}}(x)$ for the TIFN \tilde{t} as follows.

Definition 12: Let $\sigma(\tilde{t}_{\alpha})$ and $\sigma(\tilde{t}_{\beta})$ be the mean values of the intervals \tilde{t}_{α} and \tilde{t}_{β} respectively i.e.

$$\sigma(\tilde{t}_{\alpha}) = [(\chi_i - \alpha)(\bar{t} - \underline{t})]/2\chi_i$$

and

$$\sigma(\tilde{t}_{\beta}) = [(\beta - \eta_i)(\bar{t} - \underline{t})]/2(1 - \eta_i).$$

Then deviation index of the membership function $\sigma_{\mu}(x)$ and the average index of the non-membership function $\sigma_{\nu}(x)$ for the TIFN \tilde{t} are defined as

$$\sigma_{\mu}(\tilde{t}) = \int_0^{\chi_i} \sigma(\tilde{t}_{\alpha}) d\alpha = \chi_i(\bar{t} - \underline{t})/4$$

$$\text{and } \sigma_{\nu}(\tilde{t}) = \int_{\eta_i}^1 \sigma(\tilde{t}_{\beta}) d\beta = (1 - \eta_i)(\bar{t} - \underline{t})/4$$

respectively. Now we will state ranking order of TIFN. In doing that one thing we should have in mind that this ranking order is not unique and it depends on the purpose concerned. Here we will define a new ranking order based on difference between A_{μ} and A_{ν} and for that purpose we will define value judgement index $\Theta_{\tilde{a}, \tilde{b}}$ as follows

Definition 13: Let \tilde{a} and \tilde{b} are two TIFN. The average indexes of membership functions are $A_{\mu}(\tilde{a})$ and $A_{\mu}(\tilde{b})$ respectively and that of the non-membership function are $A_{\nu}(\tilde{a})$ and $A_{\nu}(\tilde{b})$ respectively. The deviation indices of membership functions are $\sigma_{\mu}(\tilde{a})$ and $\sigma_{\mu}(\tilde{b})$ respectively and that of the non-membership function are $\sigma_{\nu}(\tilde{a})$ and $\sigma_{\nu}(\tilde{b})$ respectively. Then

$$\Theta_{\tilde{a}, \tilde{b}} = \frac{\pi(\tilde{b})\{A_{\nu}(\tilde{a}) - A_{\mu}(\tilde{a})\} - \pi(\tilde{a})\{A_{\nu}(\tilde{b}) - A_{\mu}(\tilde{b})\}}{\pi(\tilde{a})\pi(\tilde{b})\max_{t=\tilde{a}, \tilde{b}}(\sigma_{\nu}(\tilde{t}) - \sigma_{\mu}(\tilde{t}))}. \quad (2)$$

Now

1) if $\Theta_{\tilde{a}, \tilde{b}} < 0$ then \tilde{a} is smaller than \tilde{b} , denoted by $\tilde{a} <_{IF} \tilde{b}$;

2) if $\Theta_{\tilde{a}, \tilde{b}} = 0$ then $\tilde{a} =_{IF} \tilde{b}$.

Here “ $<_{IF}$ ” in intuitionistic version is equivalent to “ $<$ ”

in real number set and has the linguistic interpretation “essentially less than”. Similarly “ $>_{IF}$ ” and “ $=_{IF}$ ” can be explained. For comparison of more than two TIFNs we use the notations “ \vee ” and “ \wedge ” as

$$1) \vee_i \tilde{a}_i = \tilde{a}_k \text{ if } \tilde{a}_k \text{ is largest among all}$$

$$\tilde{a}_i, i = 1, 2, \dots, n;$$

$$2) \wedge_i \tilde{a}_i = \tilde{a}_l \text{ if } \tilde{a}_l \text{ is smallest among all}$$

$$\tilde{a}_i, i = 1, 2, \dots, n.$$

Different ranking methods [16,17] have been adopted so far considering the membership and non membership function as triangular, trapezoidal or other forms of fuzzy numbers. But it is of no use when we consider the membership and non-membership functions as acceptance and rejection degree of choice of a particular thing. In this case score function is very useful. It can be defined as follows:

Chen and Tan [18] first defined a score function S_{ij} as deviation of a membership function μ_{ij} from non-membership function ν_{ij} as

$$S_{ij} = \mu_{ij} - \nu_{ij}. \tag{3}$$

Here bigger the value of S_{ij} represents bigger IFN but when S_{ij} of two IFN are same then this definition does not work. So, analyzing the deficiency of this score function Hong and Chi [19] have given a precise function as

$$H_{ij} = \mu_{ij} + \nu_{ij}. \tag{4}$$

Here also bigger the value of H_{ij} gives bigger IFN. Now these two scoring functions defined above have fundamental deficiency that they do not involve the uncertainty function π_{ij} and this seems to be very unrealistic. Liu [20] analyzing the hesitancy degree π modified the definition as

$$S1_{ij} = (\mu_{ij} - \nu_{ij})(1 + \pi_{ij}). \tag{5}$$

Now here we will use a very simple score function which is defined as

$$F(\tilde{t}) = \frac{A_\nu(\tilde{t}) - A_\mu(\tilde{t})}{\pi(\tilde{t})}. \tag{6}$$

Here one thing can be observed that

$$\Theta_{\tilde{a}, \tilde{b}} < 0 \text{ iff } F(\tilde{a}) < F(\tilde{b}) \tag{7}$$

and two properties are given as

$$1) F(\tilde{t} + \tilde{s}) = F(\tilde{t}) + F(\tilde{s});$$

$$2) F(\lambda \tilde{t}) = \lambda F(\tilde{t}), \lambda \text{ is any real.}$$

3. TIFN Matrix Game

The table showing how payments should be made at the

end of the game is called a pay-off matrix. If the player A has m strategies available to him and the player B has n strategies available to him, then the pay-off for various strategies is represented by $m \times n$ pay-off matrix. Here we consider the pay-off as TIFN

$\tilde{t} = \langle (\underline{t}, t, \bar{t}) \chi_{\tilde{t}}, \eta_{\tilde{t}} \rangle$, written in the matrix form as

$$\begin{matrix} & B_1 & B_2 & \dots & B_m \\ A_1 & \left(\begin{matrix} \tilde{t}_{11} & \tilde{t}_{12} & \dots & \tilde{t}_{1n} \end{matrix} \right) \\ A_2 & \left(\begin{matrix} \tilde{t}_{21} & \tilde{t}_{22} & \dots & \tilde{t}_{2n} \end{matrix} \right) \\ \vdots & \left(\begin{matrix} \vdots & \vdots & \ddots & \vdots \end{matrix} \right) \\ A_m & \left(\begin{matrix} \tilde{t}_{m1} & \tilde{t}_{m2} & \dots & \tilde{t}_{mn} \end{matrix} \right) \end{matrix}.$$

Here it is assumed that when player A chooses the strategy A_i and the player B selects strategy B_j it results in a pay-off \tilde{t}_{ij} to the player A .

Pure Strategy

Pure strategy is a decision making rule in which one particular course of action is selected. For fuzzy games the min-max principle is described by Nishizaki [2]. The course of the fuzzy game is determined by the desire of A to maximize his gain and that of restrict his loss to a minimum. Now for TIFN game,

$$\max - \min = \vee_i \left\{ \wedge_j \left\{ \left\langle (\underline{t}_{ij}, t_{ij}, \bar{t}_{ij}); \chi_{\tilde{t}_{ij}}, \eta_{\tilde{t}_{ij}} \right\rangle \right\} \right\}; \tag{8}$$

$$\min - \max = \wedge_j \left\{ \vee_i \left\{ \left\langle (\underline{t}_{ij}, t_{ij}, \bar{t}_{ij}); \chi_{\tilde{t}_{ij}}, \eta_{\tilde{t}_{ij}} \right\rangle \right\} \right\}. \tag{9}$$

Based on TIFN order, for such games, we define the concepts of min-max equilibrium strategies.

Definition 14 (Saddle Point): The concept of saddle point in classical form is introduced by Neumann [21]. The (k, r) th position of the pay-off matrix will be called a saddle point, if and only if,

$$\begin{aligned} \left\langle (\underline{t}_{kr}, t_{kr}, \bar{t}_{kr}); \chi_{\tilde{t}_{kr}}, \eta_{\tilde{t}_{kr}} \right\rangle &= \vee_i \left\{ \wedge_j \left\{ \left\langle (\underline{t}_{ij}, t_{ij}, \bar{t}_{ij}); \chi_{\tilde{t}_{ij}}, \eta_{\tilde{t}_{ij}} \right\rangle \right\} \right\} \\ &= \wedge_j \left\{ \vee_i \left\{ \left\langle (\underline{t}_{ij}, t_{ij}, \bar{t}_{ij}); \chi_{\tilde{t}_{ij}}, \eta_{\tilde{t}_{ij}} \right\rangle \right\} \right\} \\ \text{or, } \tilde{t}_{kr} &= \vee_i \left\{ \wedge_j \tilde{t}_{ij} \right\} = \wedge_j \left\{ \vee_i \tilde{t}_{ij} \right\}. \end{aligned} \tag{10}$$

We call the position (k, r) of entry a saddle point, the entry itself $\left\langle (\underline{t}_{kr}, t_{kr}, \bar{t}_{kr}); \chi_{\tilde{t}_{kr}}, \eta_{\tilde{t}_{kr}} \right\rangle$ the value of the game (denoted by \tilde{V}) and the pair of pure strategies leading to it are optimal pure strategies.

In Nan, Li and Zhang [13] the solution method given, involve some deficiencies which can be obviated when we use the concept, given in this paper. In [13] a reasonable solution is obtained and using it, maxmin strategy and minmax strategy are defined. But maxi-

zing the maxmin strategy and minimizing the minmax strategy does not ensure the optimality. For example, let us consider the matrix

$$\begin{matrix} & B_1 & B_2 \\ A_1 & \langle (175,180,190); 0.6, 0.2 \rangle & \langle (150,156,158); 0.6, 0.1 \rangle \\ A_2 & \langle (180,90,100); 0.9, 0.1 \rangle & \langle (175,180,190); 0.6, 0.2 \rangle \end{matrix}$$

According to the solution method defined in [13] the value of the game is

$$V_1 = \langle (152.37, 1158.44, 165.18); 0.6, 0.2 \rangle$$

although according to the method described in this paper this matrix has a saddle point (1,1) and value of the game is $V_2 = \langle (175,180,190); 0.6, 0.2 \rangle$. If we consider the comparison method of two triangular intuitionistic fuzzy numbers described in [13] we will see that $V_1 < V_2$. Hence we have got better result.

Definition 15: (TIFN expected pay-off): If the mixed strategies $x = (x_1, x_2, \dots, x_m)$ and $y = (y_1, y_2, \dots, y_n)$ are proposed by players A and B respectively, then the expected pay-off of the player A by player B is defined by

$$\tilde{E}(x, y) = \sum_{j=1}^n \sum_{i=1}^m \langle (t_{ij}, t_{ij}, \bar{t}_{ij}) \chi_{t_{ij}}, \eta_{\bar{t}_{ij}} \rangle x_i y_j. \quad (11)$$

Addition and other composition rules on TIFN which we have discussed in Definition 7 are used in this definition of expected pay-off (11). In such a situation, player A chooses x so as to maximize his expectation and player B chooses y so as to minimize player A 's maximum expectation and mathematically we write

$$\min_y \max_x \tilde{E}(x, y) = \tilde{E}(x^*, y^*) = \max_x \min_y \tilde{E}(x, y) \quad (12)$$

where (x^*, y^*) is called strategic saddle point of the game and $\tilde{V} = \tilde{E}(x^*, y^*)$ is the value of the game.

Theorem 1: If a pay-off matrix with elements as TIFN has saddle point (k, r) and \tilde{t}_{kr} is the value of the game then the pay-off matrix obtained after defuzzification with the help of score function F has also saddle point (k, r) and $F(\tilde{t}_{kr})$ is the value of the game.

Proof: If (k, r) be the saddle point of the pay-off matrix and \tilde{t}_{kr} is the value of the game then

$$\tilde{t}_{kr} = \bigvee_i \left\{ \bigwedge_j \tilde{t}_{ij} \right\} = \bigwedge_j \left\{ \bigvee_i \tilde{t}_{ij} \right\}.$$

Now using the Equations (6) and (1) we have

$$\begin{aligned} F(\tilde{t}_{kr}) &= F\left(\bigvee_i \left\{ \bigwedge_j \tilde{t}_{ij} \right\}\right) = F\left(\bigwedge_j \left\{ \bigvee_i \tilde{t}_{ij} \right\}\right) \\ \Rightarrow F(\tilde{t}_{kr}) &= \bigvee_i F\left(\bigwedge_j \tilde{t}_{ij}\right) = \bigwedge_j F\left(\bigvee_i \tilde{t}_{ij}\right) \end{aligned}$$

$$\Rightarrow F(\tilde{t}_{kr}) = \bigvee_i \left\{ \bigwedge_j F(\tilde{t}_{ij}) \right\} = \bigwedge_j \left\{ \bigvee_i F(\tilde{t}_{ij}) \right\}.$$

Therefore, (k, r) is also the saddle point of the defuzzified pay-off matrix. $F(\tilde{t}_{kr})$ is the value of the game. Hence the theorem.

Theorem 2: If (x^*, y^*) be the strategic solution of the pay-off matrix with mixed strategies then (x^*, y^*) is also the solution of the pay-off matrix after defuzzification by score function F .

Proof: Let (x^*, y^*) be the solution of the pay-off matrix then

$$\begin{aligned} &\min_y \max_x \tilde{E}(x, y) \\ &= \tilde{E}(x^*, y^*) = \sum_{i=1}^m \sum_{j=1}^n \tilde{t}_{ij} x_{ij}^* y_{ij}^* = \max_x \min_y \tilde{E}(x, y) \\ &\Rightarrow F\left(\min_y \max_x \tilde{E}(x, y)\right) = F\left(\tilde{E}(x^*, y^*)\right) \\ &= F\left(\sum_{i=1}^m \sum_{j=1}^n \tilde{t}_{ij} x_{ij}^* y_{ij}^*\right) = F\left(\max_x \min_y \tilde{E}(x, y)\right) \\ &\min_y \max_x F(\tilde{E}(x, y)) = F(\tilde{E}(x^*, y^*)) \\ &= \sum_{i=1}^m \sum_{j=1}^n F(\tilde{t}_{ij}) x_{ij}^* y_{ij}^* = \tilde{E}(F) = \max_x \min_y F(\tilde{E}(x, y)). \end{aligned}$$

Therefore, (x^*, y^*) is also a strategic solution of the defuzzified pay-off matrix and value of the game is

$$\tilde{V}(x^*, y^*) = \sum_{i=1}^m \sum_{j=1}^n F(\tilde{t}_{ij}) x_{ij}^* y_{ij}^*.$$

Hence the theorem.

4. An Application to Voting Share Problem

Suppose that there is an election where two major political parties A and B take part and total number of voters in that region is constant. It means that the increase in percentage of voters for one political party results in the same for the other political party. Suppose A has two strategies as

A_1 : Giving importance in door to door campaigning and carrying their ideology and issues to people.

A_2 : Co-operating with other small political parties to reduce secured votes of the opposition.

At the same time B takes two strategies:

B_1 : campaigning by celebrities and big rallies.

B_2 : Making lot of promises to the people.

Now the chief voting agents can not say exactly about the voting percentage but they have a certain confidence level. Still there is some hesitancy in that confidence level due to bad weather forecast. In such win-win situation we may consider the pay-offs as TIFN and the matrix is given as

$$\begin{matrix} & B_1 & B_2 \\ A_1 & \langle (4,6,9); 0.5, 0.3 \rangle & \langle (5,7,8); 0.6, 0.2 \rangle \\ A_2 & \langle (4,7,8); 0.4, 0.3 \rangle & \langle (3,5,6); 0.5, 0.2 \rangle \end{matrix}$$

Here $\langle(4, 6, 9); 0.5, 0.3\rangle$ represents that when A plays the strategy A_1 and B plays the strategy B_1 then resulting expected votes in favor of A is approximately 6 lakhs with lower bound of 4 lakhs and upper bound of 9 lakhs. The maximum confidence level and minimum non-confidence level of the Chief election agent of A are 0.5 and 0.3 respectively.

Now, let us consider the elements of the pay-off matrix as \tilde{t}_{ij} . Then using the ranking order we can get **Table 1**. Using Equation (5) we get the crisp matrix as

$$\begin{matrix} & B_1 & B_2 \\ A_1 & (6.25 & 6.75) \\ A_2 & (6.5 & 4.75) \end{matrix}$$

Since $\bigvee_i \left\{ \bigwedge_j a_{ij} \right\} = (1, 1) \neq (2, 1) = \bigwedge_j \left\{ \bigvee_i a_{ij} \right\}$, saddle point

does not exist. If we would use $\Theta_{\bar{a}, \bar{b}}$ for comparing the TIFN of original pay-off matrix we would get the same result. So, using the mixed strategy method for crisp pay-off matrix we get $x_1^* = \frac{7}{9}$, $x_2^* = \frac{2}{9}$ which are the probabilities with which player A plays the strategies A_1 and A_2 . Similarly, $y_1^* = \frac{8}{9}$ and $y_2^* = \frac{1}{9}$ are the probabilities with which player B plays with strategies B_1 and B_2 . The value of the game is 6.31 which gives optimal score in favor of player A . Since we would get same optimal strategies for player A and player B if we would use the original pay-off matrix as evident from Theorem 2, the value of the game as a TIFN is given as

$$\begin{aligned} & \frac{7}{9} \langle(4, 6, 9); 0.5, 0.3\rangle + \frac{2}{9} \langle(4, 7, 8); 0.4, 0.3\rangle \\ & = \left\langle \left(4, \frac{56}{9}, \frac{79}{9} \right); \frac{8}{90}, \frac{7}{30} \right\rangle \end{aligned}$$

which actually represents that the expected optimal votes for player A is $\frac{56}{9}$ lakhs which could be as low as 4 lakh and reach as high as $\frac{79}{9}$ lakhs. The maximum confidence level and minimum non-confidence level for the decision maker are $\frac{8}{90}$ and $\frac{7}{30}$ respectively.

Results and Discussion

This result actually represents that the expected optimal votes for player A is $\frac{56}{9}$ lakhs which could be as low as 4 lakh and reach as high as $\frac{79}{9}$ lakhs. The maximum confidence level and minimum non-confidence level for

Table 1. The computation results.

| i, j | $A_\mu(\tilde{t}_{ij})$ | $A_\nu(\tilde{t}_{ij})$ | $\sigma_\mu(\tilde{t}_{ij})$ | $\sigma_\nu(\tilde{t}_{ij})$ |
|----------------|-------------------------|-------------------------|------------------------------|------------------------------|
| $i = 1, j = 1$ | 3.125 | 4.375 | 0.625 | 0.875 |
| $i = 1, j = 2$ | 4.05 | 5.40 | 0.45 | 0.60 |
| $i = 2, j = 1$ | 2.60 | 4.55 | 0.4 | 0.7 |
| $i = 2, j = 2$ | 2.375 | 3.8 | 0.375 | 0.6 |

the decision maker are $\frac{8}{90}$ and $\frac{7}{30}$ respectively. Unlike [13], in this paper, we didn't go for the reasonable solution and instead, we tried to reach to the optimality with the help of a crisp pay off matrix. The reasonable solution [13] does not confirm the optimality but the theorems 1 and 2 support that the optimality exist when we use this method.

5. Conclusion

In this paper, we have used TIFN as elements of pay-off matrix. As a result we have considered here players' preference information in terms of support, opposition and neutralization and also his confidence and non-confidence level about the approximation. Using the definitions and operations of TIFN we have described a ranking order based on the definition of value judgement index. Then we have described score function to defuzzify the matrix game and made a comparative study on scoring function approach and an IF approach in voting share problem. The merit of this methodology is that it obtains a deterministic solution of a matrix game with IF pay-off. There is a scope to apply such a methodology in other conflicting interest problems.

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