

Energy Aware Clustered Based Multipath Routing in Mobile Ad Hoc Networks

M. BHEEMALINGAIAH¹, M. M. NAIDU², D. Sreenivasa RAO³

¹*Dept. of Computer Science and Engineering J.N.T University, Hyderabad, India*

²*Dept. of Computer Science and Engineering, S.V.U College of Engineering, Tirupati, India*

³*Dept. of Electronics and Communication Engineering, J.N.T University, Hyderabad, India*

Email: saibheem2008@gmail.com, mmnaidu@yahoo.com, dsraoec@yahoo.co.uk

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Abstract

With the advance of wireless communication technologies, small-size and high-performance computing and communication devices are increasingly used in daily life. After the success of second generation mobile system, more interest was started in wireless communications. A Mobile Ad hoc Network (MANET) is a wireless network without any fixed infrastructure or centralized control; it contains mobile nodes that are connected dynamically in an arbitrary manner. The Mobile Ad hoc Networks are essentially suitable when infrastructure is not present or difficult or costly to setup or when network setup is to be done quickly within a short period, they are very attractive for tactical communication in the military and rescue missions. They are also expected to play an important role in the civilian for as convention centers, conferences, and electronic classrooms. The clustering is an important research area in mobile ad hoc networks because it improves the performance of flexibility and scalability when network size is huge with high mobility. All mobile nodes operate on battery power; hence, the power consumption becomes an important issue in Mobile Ad hoc Network. In this article we proposed an Energy Aware Clustered-Based Multipath Routing (EACMR), which forms several clusters, finds energy aware node-disjoint multiple routes from a source to destination and increases the network life time by using optimal routes.

Keywords: Clustering, CONID, MANET, Mutlipath, AODVM

1. Introduction

The history of wireless networks started in the 1970s and the interest has been growing ever since. Based on infrastructure, the wireless networks broadly classified into two types, first type infrastructure networks contains base-stations; an example of this wireless network is the cellular-phone network where a phone connects to the base-station with the best signal quality. When the phone moves out of range of a base-station, it does a "hand-off" and switches to a new base-station within reach, the second type is called as Mobile Ad hoc Network without any fixed infrastructure or centralized control; it contains mobile nodes that are connected in an arbitrary manner. It enables the users to communicate without any physical infrastructure regardless of their geographical location. All nodes in the Mobile Ad hoc Network (MANET) are

mobile and can be connected dynamically in an arbitrary manner. Each node behaves as a router, takes part in discovery and maintains the routes to other nodes. The nodes are the main components of the network; these nodes can move freely at any time and can leave or join the network so the network structure changes dynamically due to mobility [1].

1.1. Clustering and Multipath Routing

In a clustering scheme the mobile nodes in a MANET are divided into different virtual groups based on certain rules. The mobile nodes may be assigned a different status such as clusterhead, clustergateway or cluster-member. A clusterhead normally serves as a local coordinator for its cluster, performing intra-cluster transmission arrangement, data forwarding, and so on. A cluster-

gateway is a non-clusterhead node with inter-cluster links, so it can access neighboring clusters and forward information between clusters. A cluster member is usually called an ordinary node [2].

A fundamental problem in the MANET is how to deliver data packets among nodes efficiently without predetermined topology or centralized control, which is the main objective of ad hoc routing protocols. Because of the dynamic nature of the network, ad hoc routing faces many unique problems not present in wired networks. Particularly in MANETs where routes become obsolete frequently because of mobility and poor wireless link quality. The Multipath routing addresses these problems by providing more than one route to a destination node. Multipath routing appears to be a promising technique for ad hoc routing protocols, the multiple paths can be useful in improving the effective bandwidth of communication, responding to congestion and heavy traffic, increasing delivery reliability and security. The traffic can be distributed among multiple routes to enhance transmission reliability, provide load balancing, and secure data transmission [3].

Most existing routing protocols for MANET build and utilize only one single route for each pair of source and destination nodes. Due to node mobility, node failures, and the dynamic characteristics of the radio channel, links in a route may become temporarily unavailable and making the route invalid. The overhead of finding alternative routes may be high and extra delay in packet delivery may be introduced. Multipath routing addresses this problem by providing more than one route to a destination node. Source and intermediate nodes can use these routes as primary and backup routes. Alternatively, source node can distribute traffic among multiple routes to enhance transmission reliability, provide load balancing, and secure data transmission.

2. Related Work

This work falls under three areas in the MANET: clustering algorithms, multipath routing protocols and energy aware routing protocols, this section reviews these areas focusing on their relationship to the proposed protocol.

Jane Y. *et al.* [2] surveyed the different clustering mechanisms and described the advantages and disadvantages of each clustering scheme. The clustering schemes broadly classified into six categories Dominating Set based clustering schemes, Low maintenance clustering schemes, Mobility aware clustering schemes, Energy efficient clustering schemes, Load balancing clustering schemes and Combined-metrics based clustering schemes. The Energy efficient clustering avoids unnecessary energy consumption or balancing energy consumption for mobile nodes in order to prolong the lifetime of mobile nodes and the network.

The on-demand routing is the most popular approach in the MANET. Instead of periodically exchanging route messages to maintain a permanent route table of the full topology, the on-demand routing protocols build routes only when a node needs to send the data packets to a destination. The standard protocols of this type are Dynamic Source Routing (DSR) [4] and the Ad hoc On-demand Distance Vector (AODV) routing [5]. However, these protocols do not support multipath.

The several multipath on-demand routing protocols were proposed, some of the standard protocols are, the Ad hoc On-demand Multipath Distance Vector (AOMDV) [6] is an extension to the AODV protocol for computing multiple loop-free and link-disjoint paths. The Split Multipath Routing (SMR) [7] is an on-demand Multipath source routing protocol can find an alternative route that is maximally disjoint from the source to the destination. The Multipath Source Routing (MSR) [8] is an extension of the DSR protocol to distribute traffic among multiple routes in a network. The Ad hoc On-demand Distance Vector Multipath Routing (AODVM) [9] is an extension to the AODV for finding multiple node disjoint paths. These protocols build multiple routes based on demand but they did not consider energy.

Several energy aware multipath on-demand routing protocols have been proposed [10–14], but these protocols are not based on clustering. The Cluster Based Multipath Dynamic Source Routing CMDSR [15] was designed to be adaptive according to network dynamics. It uses the hierarchy to perform route discovery and distributes traffic among diverse multiple paths, but it does not consider energy in the route selection.

The Grid-Based Energy Aware Node-Disjoint Multipath Routing Algorithm (GEANDMRA) [16] considers energy aware and node-disjoint multipath, it uses grid head election algorithm to select the grid-head which is responsible for forwarding routing information and transmitting data packets. The routing is performed in a grid-by-grid manner, the network area is partitioned into non-overlapping square zones with the same size, each zone is named with a unique ID(x,y) as the conventional coordinate. At any time each node can obtain its location information from Global Position System (GPS) to know in which zone it is located. But the disadvantages of GPS are as follows [17].

- There are several factors that introduce error to GPS position calculations. A major source of error arises from the fact that radio signal speed is constant only in a vacuum. It means that distance measurements may vary as the values of the signal speed vary in the atmosphere. Water vapor and other particles in the atmosphere can slow signals down, resulting in *propagation delay*.
- Another source of errors due to *multipath fading*, which occurs when a signal bounces off a building

or terrain before reaching the receiver's antenna, also can reduce accuracy.

- Another factor affecting the precision is satellite geometry.
- The largest source of potential error is *selective availability*, an intentional degradation of L1, the civilian GPS signal. SA was originally intended to prevent a hostile force or terrorist group from exploiting the technology.
- An location fix can only be identified if a 'fix' from at least 3 satellites are available.
- In addition, distance measurements are less reliable when the receiver of satellite locks on to or closely oriented with respect to each other. Atomic clock discrepancies, receiver noise, and interruptions to ephemeris monitoring can result in minor errors.
- GPS does not work effectively if the 'line of sight' between a device and the satellites is obscured. Not only does this render GPS ineffective indoors but also if any material gets in the way of this 'line of sight'. It is fair to assume therefore that a worker in a vehicle or in a station or operating in a dense urban area with tall buildings may struggle to get a location fix from GPS.
- The device needs a dedicated GPS antenna, chipset and software, so there is a additional cost for this hardware and software.
- Power consumption in the device is significant for using GPS.

We proposed novel approach without using GPS; this approach is very suitable and applicable for the cluster based mobile ad hoc networks.

3. Energy Aware Clustered-Based Multipath Routing (EACMR)

The HELLO messages are already in usage in on demand routing, each node maintains only the status of its neighbors. Periodically, each node broadcasts a hello message to all its neighbors to indicate its active status. Each node receives the hello messages from its neighbors and updates its neighbor information table. In the EACMR, the clusters are formed by using hello messages, this result in less overhead. The EACMR finds node-disjoint multiple routes from a source to destination and increases the network life time by using optimal routes.

3.1. Network Model

A MANET is represented by undirected graph, $G=(V,E)$ where V is the set of nodes and E is the set of bidirectional links. It is assumed that, at any time, nodes are

situated randomly throughout the network area in accordance with a two-dimensional uniform random variable distribution. Each node is equipped with a single network interface card (NIC) and has a transmission radius of r . Each node has mobility, the speed is uniformly chosen between the minimum and maximum speeds. When the node reaches to its destination, it stays there for a certain pause time, after which it chooses another random destination point and repeats the process, the mobility is defined as the distance moved per unit time by a node in the network. Suppose at time t_1 the node n_i is at (x_1, y_1) and by time t_2 the node n_i has moved to (x_2, y_2) , then the mobility of the node n_i denoted by

$$M_{n_i} = \frac{1}{(t_2 - t_1)} \sqrt{(x_2 - x_1)^2 + (y_2 - y_1)^2}$$

All mobile node relay on battery, the energy consumption varies from 240mA at receiving mode and 280mA in the transmitting mode using 0.5V energy. Thus,when calculating the energy consumed to transmit a packet p is $E_{tx}(p)=I \times V \times t_p$. Joules are needed [18], here, I is the current, V is the voltage and t_p is the time taken to transmit the packet p .

The energy required to transmit a packet p is given by $E_{tx}(p)=280mA \times V \times t_p$. The energy is required to receive a packet p is given by $E_{rx}(p)=240mA \times V \times t_p$. The energy consumption of overhearing the data transmission may be assumed as equivalent to energy consumption of receiving of the packet. When a packet is transmitted by n_i at time t then it updates its residual battery capacity by using the following formula $R_i^t = R_i^t - E_{tx}(p)$. when packet is received by n_i at time t then it updates its residual battery capacity by using the following formula $R_i^t = R_i^t + E_{rx}(p)$.

3.2. Cluster Formation and Maintenance

The Combined Higher Connectivity Lower ID (CONID) clustering algorithm [19] is used generates the clusters in the network. It is an extension of the lowest ID algorithm; the lowest ID algorithm does not take into account the connectivity (degree) of nodes, and therefore may produce more number of clusters than necessary. The pure connectivity based clustering algorithm modified version the lowest ID algorithm, in which ID is replaced by node degree, but it does not work properly because of numerous ties between nodes. The CONID uses node degree as the primary key and ID as the secondary key in cluster decisions. It is generalized the connectivity to count all k -hop neighbors of the given node. For $k=1$, the connectivity is equivalent to node degree, whenever the connectivities are the same, IDs are compared to make the decision. The clustering algorithm, referred to as the k -CONID (k -hop connectivity ID) algorithm, works as follows.

Each node is assigned with clusterhead priority. A pair is denoted by $did=(d, ID)$, where d is its connectivity and ID is its IP address.

Let $did'=(d', ID')$ and $did''=(d'', ID'')$. Then $did' > did''$ if $d' > d''$ or $d'=d''$ and $ID' < ID''$. That is, a node has clusterhead priority over the other node if it has higher connectivity or in case of equal connectivity and has lower ID. The cluster formation and cluster maintenance are clearly explained in [19], after applying clustering algorithm the network is shown Figure 1. After running CONID, each clusterhead finds its all neighbor clusterheads, for example, clusterheads 18 and 10 are the neighbor clusterheads of clusterhead 2 in Figure 1.

3.3. Route Selection

The route selection is based on cost function; the main objective is to give more weight (or) cost to node with less energy to prolong its life time. Let R_i^t be the battery capacity of a node n_i at time t . Let $f_i(R_i^t)$ be the cost function of node n_i , it is also considered as node's cost or weight of node n_i at time t , the cost of node n_i at time t is inversely proportional to its residual energy i.e. $f_i(R_i^t) = 1/R_i^t$. As the battery capacity decreases, the value of cost for node n_i will increase. The main objective of route selection is to select an optimal path based on costs of clusterheads, because clusterhead normally serves as a local coordinator for its cluster, performing intra-cluster transmission arrangement, data forwarding. So clusterhead is important node within cluster and spends its energy for other nodes.

We define the following cost function to a clusterhead. Let C_i be i th cluster whose clusterhead at time t is denoted by CH_i . The Cost of CH_i at time t as follows

$$Cost(CH_i) = \rho_i \times \left[\frac{F_i}{R_i^t} \right] \times w_i \times N(C_i)$$

where

- ρ_i : Transmit power of CH_i
- F_i : Full-charge capacity of CH_i
- R_i^t : Remaining battery capacity of CH_i at time t .
- w_i : weight factor of CH_i , which depends upon various factors, like battery's quality, battery's capacity, life time, battery's back up, price.

$N(C_i)$: the size of cluster C_i , it is the total number of all the nodes (clustermembers, gateways and clusterhead) in C_i , it is directly proportional to the cost of the node.

Let P_j be the path from source node S to destination node D via clusterheads $CH_1, CH_2, CH_3, \dots, CH_n$ at time t , it is denoted by

$$P_j = S - CH_1 - -CH_2 - -CH_3 - - - - -CH_n - D$$

On the above path, in between two clusterheads, either gateway or clustermember (non-clusterhead nodes) may exist, however they are not consider in the evaluation of

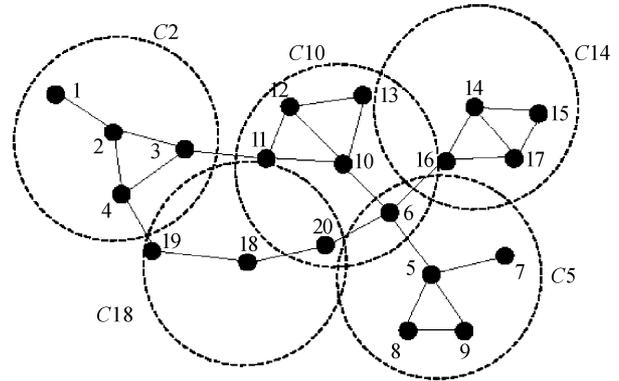


Figure 1. Network with clusters.

cost of the path. Because main intension is to consider the total cost of path through clusterheads only, we define the cost of path P_j at time t via clusterheads is the sum of the costs of all the clusterheads on P_j , it is denoted by

$$Cost(P_j) = \sum_{i=1}^n Cost(CH_i)$$

3.3.1 Optimal Path Selection

Let k be number of node disjoint paths from source S to destination D , an optimal path is a path whose cost is least among k paths.

The optimal path is also called primary path. Initially it is used for data transmission, if it fails, then the secondary path (whose cost is least among $k-1$ paths) is used for data transmission. For example in figure 2, there are three node disjoint paths from source S to destination D via clusterheads; the paths are 1-2-3-4-5, 1-7-6-5 and 1-8-9-10-5.

If source node S belongs to i th cluster whose clusterhead is CH_i then it is called as source clusterhead. Similarly, If destination D belongs to j th cluster whose clusterhead is CH_j then it is called as destination clusterhead in Figure 2. Clusterhead 1 is called as source clusterhead and clusterhead 5 is called as destination clusterhead. If

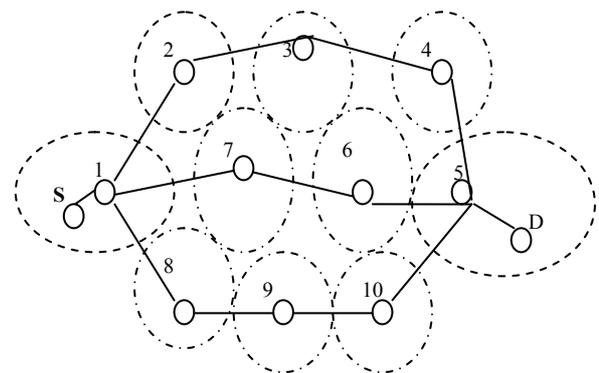


Figure 2. Node-disjoint paths through Clusterheads.

two clusterheads are neighbors then they are called as 1-hop neighbor clusterheads. In Figure 2, the clusterheads 3, 7, 8 are neighbors of the clusterhead 1.

3.4. Route Discovery

The route selection is based on the route discovery. In order to facilitate the computation of multiple node disjoint paths from the source to destination, We choose the Ad hoc On-Demand Distance Vector Multipath (AODVM) [9] protocol as a candidate protocol and make modifications to it to enable the discovery of node disjoint paths via clusterheads. The AODVM is the extension of AODV [6] to provide multiple nodes disjoint paths.

3.4.1. The Proposed Modifications

The proposed modifications are explained briefly as follows. Only clusterheads maintain routing tables and run this protocol to find nod-disjoint paths. The other nodes (Gateway nodes and Cluster members) don't maintain the routing tables, simply they forward the packets according to specified path.

3.4.2. Modification of Control Packets

The RREQ packet of AODVM is same as RREQ packet of EACMR. The RREP packet of AODVM is extended as RREP packet of EACMR by adding with *Cost field*, this field carries cumulative costs of clusterheads through which it passes. The initial value of this field is zero.

3.4.3. Modification of Tables

In the AODVM, each node maintains two tables, the routing table is used to forward the data packets from source to destination where as the RREQ table is used to form the route from source to destination, the fields of both tables are shown in Figure 3 and Figure 4.

In the EACMR, the routing table is extended to include the cost field. But RREQ table is same. These two tables are maintained by clusterheads only.

3.4.4. The Proposed Modifications in Node Functioning

In the AODVM, when the source node wants to send packets to a destination, it looks up its route cache to determine if it already contains a route to the destination. If it finds that an unexpired route to the destination exists, then it uses this route to send the packet. But if the node

does not have such a route, then it initiates the route discovery process by broadcasting a RREQ packet.

It is modified that when the source node wants to send packets to a destination, it sends the RREQ to its clusterhead; clusterhead looks up its route cache (routing table) to determine if it already contains a route to the destination. If it finds that an unexpired route exists to the destination, then it sends the reply to the source and uses this route to forward the packets from the source node, else then it initiates the route discovery process by forwarding a RREQ packet to its all neighbor clusterheads.

In the EACMR, the propagation of RREQ packet at intermediate clusterhead follows the same rules as at intermediate node in AODVM. When RREQ packet arrives at intermediate clusterhead, The RREQ is inspected and the following two cases it is dropped.

- 1) If this RREQ packet was already processed by the intermediate clusterhead or
- 2) If TTL value of the RREQ packet becomes zero.

Otherwise, the intermediate clusterhead broadcasts the RREQ to all its neighbor clusterheads by recoding the information about the clusterhead through which RREQ was received and corresponding hop-count back to the source clusterhead into its RREQ table, in this way the RREQ is forwarded by a intermediate clusterheads. Finally several RREQs will reach to destination clusterhead through different paths from the source clusterhead.

When a RREQ packet is received to a destination clusterhead, it generates a RREP packet and sent back to its neighbor clusterhead (last hop) from which the RREQ packet has been received. When an intermediate clusterhead receives the RREP packet, it checks its RREQ table to find its neighbor clusterhead (next hop in reverse path) through which shortest path back to the source clusterhead and sends the RREP packet to it by deleting corresponding entry in the RREQ table.

In order to ensure that a clusterhead does not participate in multiple paths, If it receives another RREP packet from different neighbor clusterhead, it is dropped (i.e., RREQ table is already empty), it generates a Route Discovery Error (RDER) packet and sends it back to the neighbor clusterhead through which another the RREP packet has been received. The neighbor clustered upon receiving the Route Discovery Error (RDER) packet will try to forward the RREP packet to another neighbor clusterhead through which shortest path back to the source clusterhead and sends the RREP packet to it by deleting corresponding entries in the RREQ table.

Finally several RREP packets will be received by the source clusterhead, this information is recorded into its cache based on arriving order and it sends the reply to the source. The path with minimum cost is selected as an optimal path for transmission of source's data. If clus-

Dest.	Source	Last hop	Next hop
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Figure 3. Fields of routing table.

Dest.	Source	Neighbor who transmitted the RREQ	hop	Expiration time
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Figure 4. Fields of RREQ table.

terhead wants to transmit the data, it follows above procedure to find node disjoint paths.

3.5. Route Maintenance

The EACMR handles route maintenance in a manner similar to the AODVM. Whenever a link breakage happens in a route due to a node moving away, the previous hop node of the moved away node is responsible for sending a Route Error (RERR) message back to the source clusterhead to inform the breakage. It chooses alternative routes to maintain the connection. If there are no more redundant routes left, then it will start a new route discovery.

3.6. Congestion Control and Increasing Network Lifetime

When a congestion state occurs in a routing path, congested clusterhead sends choke packet to the source clusterhead, it distributes the incoming data packets to the other node-disjoint routing paths to avoid the congestion.

During the transmission of data packets, if battery energy of clusterhead is reached to threshold energy, then it selects a cluster member as new clusterhead whose energy is high among its cluster members, then it sends information about cluster and it will act as cluster member.

4. Simulation

In this section, it describes the simulation, various chosen parameters for simulation and the various performance metrics. With reference to simulation model in [20], we have designed own simulator with the following modules.

Network Formation Module: This module is used to generate a random network, inputs of this module are simulation area (length x breadth), number of nodes, cell radius of each node, initial position of each node and initial energy of each node. The output of this module is a random network.

Node Mobility Module: This module sets the speed, direction and pause time of each node and allows each node to move in random direction. All the nodes in an ad hoc network are mobile. In this simulation, Random waypoint mobility model is used with pause time of each node is 10 sec and speed of each mobile node is 0 to 2m/sec.

Route Requests Event Generator Module: This module accepts the number of route requests from user, and then selects source and destination pairs randomly. Each route request follows the poison distribution process and each call duration time follows exponential distribution.

EACMR module: This is core module that incorporates several functions like Cluster formation, cluster maintenance, route discovery, route selection, route maintenance and congestion control.

Computation module: This module estimates various performance metrics like number of clusters, number of border nodes, the power consumption, residual energy, number of nodes expired (reached to threshold), overhead, throughput, end-to-end delay and other parameters.

Considering fifty nodes randomly distributed in an area of 1000m x 1000 m. It is assumed that the channel bandwidth is 2Mbps, a free space radio propagation model in which the signal power attenuates is $1/r^2$, where r is the distance between the nodes. Each node is equipped with a single network interface card and has a transmission radius of $r=14m$. The distributed coordination function of IEEE 802.11 is assumed at MAC layer.

All nodes operate in promiscuous mode, so it can overhear packets destined for others. It is assumed that the transmission power, receiving power are fixed for all the nodes and two nodes can hear each other if their distance is in the transmission range. The speeds are uniformly chosen between the minimum and maximum speeds and are set to 0m/s and 2m/s, respectively. When the node reaches its destination, it stays there for a certain pause time, after which it chooses another random destination point and repeats the process. The simulation ends after 100s. It is assumed that transmission ranges of all the nodes are equal. All nodes are assumed to have the same amount of battery capacity with full energy at the beginning of the simulation. Initial energy of each node is set to 100 Joules. The equal weight factor is chosen for all clusterheads. For the mobile scenarios, the random waypoint model is used to node mobility. In this model, a node chooses a random point in the network, and moves towards that point at a constant speed.

It is assumed that the number of route requests is denoted by λ and follows the poison process where as call holding time follows exponential distribution. When a route request occurs, two nodes are randomly selected as source and destination. The data traffic is generated by Constant Bit Rate (CBR) sessions initiated between the source and destination. Each clusterhead maintains threshold value (cut-off). The table 1 shows simulation parameters and their values.

During simulation, several performance parameters were estimated, like number of clusters (NC, number of border nodes, CH%, BP%, ASP, the total energy consumption at each clusterhead and total residual energy at each clusterhead. Along the number of clusterheads reached to the threshold, the results are shown in Table 2.

Table 1. Simulation parameters.

Sno	Parameter	Value
1	Simulation area and Network Size	1000mx1000m, 50 Nodes
2	Transmission Range	88 meters
3	Transmission Power	0.7 Joule/packet,
	Receiving Power	0.3 Joule/packet
4	Node Mobility Model, Pause Time and Speed	Random waypoint mobility model, 10 sec, 0 to 2m/s
5	Initial Energy, Maximum Battery Capacity	100 Joules, 100 Joules
6	Weight factor of clusterhead	1
7	Threshold Value	5 Joules
8	Route request arrival rate λ	5, 10 per 10 sec
9	Traffic type, Maximum Data Packet size	Constant Bit Rate (CBR), 512 bytes
10	Queue type and queue Size	Drop tail, 60*512 bytes
11	Total Simulation Time	100 sec

Table 2. Results.

Sno	Parameter	D=3.02	D=3.24	D=3.8
1	Number of Clusters	13	13	13
2	Elected Clusterheads	2, 19, 6, 42, 40, 33, 28, 26, 22, 37, 15, 12, 48	2, 18, 5, 40, 32, 28, 24, 36, 50, 42, 14, 46, 10	2, 19, 6, 42, 40, 32, 28, 24, 36, 50, 17, 12, 46
3	Number of Border nodes	9	11	13
4	CH%	0.26	0.26	0.26
5	BP%	0.18	0.22	0.26
6	AS	3.846	3.846	3.846

The size and degree are denoted by N and D, respectively. CH% denotes the ratio of clusterhead nodes (that is, the number of clusters divided by the total number of nodes). Border nodes are nodes that belong to more than one cluster (that is, which are at distance at most k hops from at least two CHs), BP% denotes the ratio of border nodes, (that is, the number of border nodes divided by the total number of nodes). AS is the average size of a cluster (the average number of nodes in a cluster, that is, the number of nodes divided by the total number of clusters). Other results are shown in Figure 5 to Figure 7.

The total power consumption is directly proportional to various factors like network size, route requests arrival rate, packet arrival rate, packet size (header size and payload size), packet collision and retransmissions. Total residual energy is indirectly proportional to the power consumption. The network life depends on then node expiration which in turn depends upon energy consumption and threshold value. The node life time is indirectly proportional to the node's energy consumption and it is also directly proportional to the threshold value of the energy of each node is denoted by γ . It is also called cut-off value; it is always greater than one. During the transmission of data, each node checks whether its energy

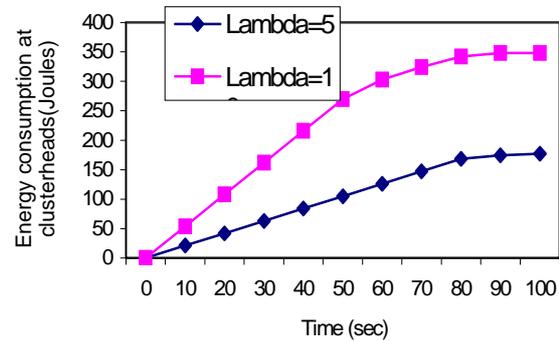


Figure 5. Energy consumption Vs time.

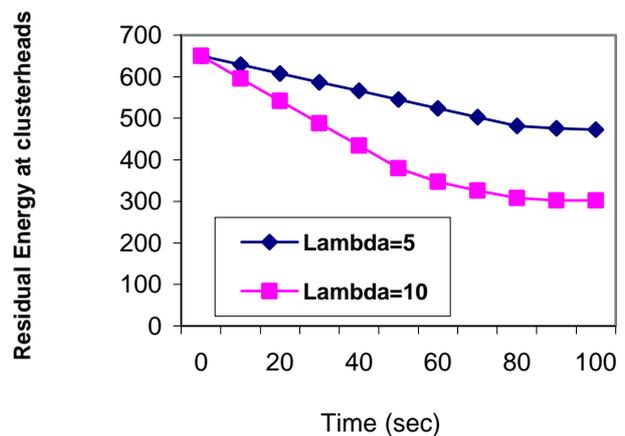


Figure 6. Residual energy Vs time.

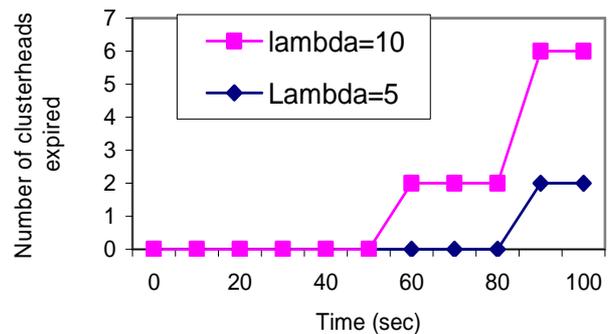


Figure 7. Number of CHs expired Vs time.

reaches to threshold or not. If its energy reaches to threshold then it will expire. The network lifetime can be defined in many ways:

- It may be defined as the time taken for K% of the nodes in a network to die
- It might be the time taken for the first node to die

- It can also be the time for all nodes in the network to die

It can also be the time for all nodes in the network to die. To maximize the lifetime of network, each clustered maintains threshold.

Figure 5 depicts the variation of total energy consumption of clusterheads versus time; Figure 6 depicts the variation of total residual energy of clusterheads versus time. Figure 7 depicts number of clusterheads expired versus time.

5. Conclusions

Proposed novel approach (EACMR) is very suitable and applicable for the cluster based mobile ad hoc networks. It is not based on GPS. It can be applied to a mobile ad hoc network that is using any clustering scheme. In this work, the CONID clustering scheme is used as background to form clusters, instead of that any clustering scheme may be used to form clusters in the network.

The EACMR is designed to find energy aware node-disjoint multiple routes from a source to a destination through clusterheads. It increases the network life time by using optimal routes, as compare to on demand multipath routing protocols, it significantly reduces the total number of route request packets using clustering technique, this result in an increased packet delivery ratio, decreasing end-to-end delays for the data packets, lower control overhead, fewer collisions of packets, decreasing power consumption. It supports the reliability by using multiple node-disjoint paths.

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