

The Effect of Price Discount on Time-Cost Trade-off Problem Using Genetic Algorithm

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Abstract

Time-cost trade off problem (TCTP), known in the literature as project crashing problem (PCP) and project speeding up problem (PSP) is a part of project management in planning phase. In this problem, determining the optimal levels of activity durations and activity costs which satisfy the project goal(s), leads to a balance between the project completion time and the project total cost. A large amount of literature has studied this problem under various behavior of cost function. But, in all of them, influence of discount has not been investigated. Hence, in this paper, TCTP would be studied considering the influence of discount on the resource price, using genetic algorithm (GA). The performance of proposed idea has been tested on a medium scale test problem and several computational experiments have been conducted to investigate the appropriate levels of proposed GA considering accuracy and computational time.

Keywords: Project Management, PERT Networks, Time Cost Trade off, Genetic Algorithm, Discount

1. Introduction

Time-cost trade off problem (TCTP) is an important part of the project management in planning phase. It is a sub problem of project scheduling when finding the most cost effective way to finish a project within time limit is desirable. In the TCTP, the objective is to determine the duration of each activity in order to obtain the minimum costs of the project. All existing patterns of this problem satisfy the objective function through expediting the activity durations. In this problem, determining the optimal levels of activity durations and activity costs leads to a balance between the project completion time and the project total cost. There are some procedures in TCTP (e.g. extra resource allocation, improvement in technology level, increase in the quality of materials, hiring highly skillful human resources, etc.) used for expediting the activity durations, according to activity characteristics. These procedures can be summarized to a unique cost function corresponding to each activity.

Herroelen and Leus [1] evaluated some stochastic

models of the TCTP under two main categories: "The stochastic discrete time/cost trade-off problem" and "Multi-mode trade-off problem in stochastic networks". In the stochastic TCTP, when objective function of the model is related to the time, Bowman [2] presents a heuristic algorithm based on simulation technique and solves a general project compression problem using the derivative estimators. Besides, the application of mathematical programming for the stochastic project crashing problem was suggested by Abbasi and Mukattash [3], in which activities are assumed to have three time estimates. Unfortunately, it seems that the proposed approach has been developed for a network with a single path and is not applicable for the networks with slight differences in the path durations. Arisawa and Elmaghraby [4] suggested the use of fractional linear programming for the optimal allocation of resources to the activities in order to maximize the reduction in project mean duration per unit investment in GERT type networks. A novel resource-constrained project scheduling problem has been modeled by Golenko and Gonik [5] as a knapsack resource reallocation problem in which a heuristic algorithm has been to determine the optimal amount of resources allocated to the activities and their start times.

Some authors have investigated the stochastic TCTP to minimize the project total costs [6–8]. Gutjahr et al. [9] studied the discrete model of this problem by using the stochastic branch and bound method to improve the probability of meeting project deadline. In their work, they assumed that the crashing of activity durations (by using additional costs) leads to higher direct costs and lower penalty costs. In another research, Mitchell and Klastorin [10] have formulated the objective function with the direct, indirect, and penalty costs. This study presents a Stochastic Compression Project (SCOP) heuristic based on decomposition of PERT networks into the serial and parallel subnets. Recently, some researchers have considered the multi criteria/objective formulation for the stochastic TCTP [11–14].

In all of the presented papers, TCTP have investigated under various behaviors of the cost function such as discrete cost function [5,15-19], linear continuous cost function [20–22], nonlinear convex cost function [23,24], nonlinear concave cost function [25] and linear piecewise cost function [26,27]. But, in all of them, influence of discount has not been investigated. In this study, we develop a new TCTP based on discount in resource price to maximize the project completion probability in a predefined deadline using a limited available budget. In order to construct the model we assume that activity durations follow the normal distribution and the activity mean durations represent the decision variables. In order to solve this problem, we organize the genetic algorithm (GA) technique. Recently, some researchers use the GA in order to solve the project scheduling problems such as TCTP [11].

This paper is organized in the following way. The mathematical model is presented in Section 2. Section 3 describes the structure of the proposed approach, which is based on the GA algorithm. The characteristics of a large scale numerical example and results of applying the proposed approach to that example are presented in Section 4. The appropriate levels of GA parameters are investigated in this section. Finally Section 5 consists of concluding remarks.

2. Mathematical Model

Let G = (N,A) be an acyclic Activity on Arrow (AOA) graph with arrow set A and node set N, where the

source and sink node are denoted by s and t, respectively. $N = (n_1, n_2, ..., n_m)$ represents the set of events, $A = (a_1, a_2, ..., a_n)$ represents the set of activities in a project PERT network with m nodes and n activities.

Considered notations are presented as follows:

(i, j) Activity with i head node and j tail node

 t_{ii} Random variable of activity (i, j) duration

 μ_{ij} Mean duration of activity (i, j) (decision variable)

 σ_{ii} Standard deviation of activity (i, j)

 CS_{ii} Cost slope of activity (i, j)

 R_{ij} Amount of allocated resource to activity (i, j)

 T_d Project completion deadline

 μ_{ii}^{u} Upper limit of mean of ith activity.

 μ_{ii}^{l} Lower limit of mean of ith activity.

M Amount of available budget (i, j).

L Total number of paths.

 n_r Number of activities which lie on path r.

 T_r Random variable of path r duration.

 μ The mean duration of path r.

 σ_r The standard deviation of path r.

φ Cumulative distribution function (CDF) of normal standard distribution.

Generally, the TCTP is a nonlinear optimization problem. In order to investigate the effect of discount on this problem, we extract the objective function (project completion probability) and constraints, separately. The objective function (OF) is concerned with the maximization of project completion probability. To construct the OF, we can write a single path completion probability according to the central limit theorem (CLT), in the following way:

$$\Pr(T_r < T_d) = \Pr\left(Z < \frac{T_d - \mu_r}{\sigma_r}\right) = \Pr\left(Z < \frac{T_d - \sum_{ij \in r} \mu_{ij}}{\sum_{ij \in r} \sigma_{ij}}\right) = \phi\left(\frac{T_d - \sum_{ij \in r} \mu_{ij}}{\sum_{ij \in r} \sigma_{ij}}\right)$$
(1)

Equation (1) is related to a single path and can't improve the total project completion probability. Thus, considering all of the paths, the following objective function is suggested:

$$Max \{ Min \ \phi \left(\frac{T_d - \sum_{ij \in r} \mu_{ij}}{\sum_{ij \in r} \sigma_{ij}} \right) \quad \forall \ r = 1, 2,, L \ \}$$
 (2)

Two type of restrictions should be modeled in a mathematical formulation to ensure the obtained solu-

tions satisfy the budget limitation and real conditions. Each activity can be planned between two maximum and minimum limit of its mean duration. Following mathematical representation show this type of constraint:

$$\mu_{ij}^{l} \le \mu_{ij} \le \mu_{ij}^{u} \tag{3}$$

It is assumed that each activity needs a unique type of resource. Using this assumption, in order to model the budget limitation, we can define the effect of discount on resource price in a mathematical representation. For this purpose, cost slope (CS) due to various levels of μ_{ij} is presented below:

Using (4) we can formulate the cost function for activity (i, j) as a mathematical representation, in the following way:

$$CS_{ij} = \begin{cases} S_{ij}^{\circ} & R_{ij}^{\min} \leq R_{ij} < R_{ij}^{1} \\ S_{ij}^{1} & R_{ij}^{1} \leq R_{ij} < R_{ij}^{2} \\ S_{ij}^{2} & R_{ij}^{2} \leq R_{ij} < R_{ij}^{3} \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ S_{ij}^{k-1} & R_{ij}^{k-1} \leq R_{ij} < R_{ij}^{k} \\ & & & & \\ S_{ij}^{k} & R_{ij}^{k} \leq R_{ij} < R_{iji}^{max} \end{cases}$$

$$(4)$$

$$C_{ij}(\mu_{ij}) = \begin{pmatrix} S_{ij}^{\circ}(\mu_{ij}^{\max} - \mu_{ij}) & \mu_{ij}^{\circ} \leq \mu_{ij} < \mu_{ij}^{\max} \\ \vdots & \vdots \\ S_{ij}^{\circ}(\mu_{ij}^{\max} - \mu_{ij}^{\circ}) + \sum_{k=1}^{n-1} S_{ij}^{k}(\mu_{ij}^{k-1} - \mu_{ij}^{k}) + S_{ij}^{n}(\mu_{ij} - \mu_{ij}^{n-1}) & \mu_{ij}^{n-1} \leq \mu_{ij} < \mu_{ij}^{n} \\ \vdots & \vdots \\ S_{ij}^{\circ}(\mu_{ij}^{\max} - \mu_{ij}^{\circ}) + \sum_{k=1}^{\min-1} S_{ij}^{k}(\mu_{ij}^{k-1} - \mu_{ij}^{k}) + S_{ij}^{\min}(\mu_{ij} - \mu_{ij}^{\min-1}) & \mu_{ij}^{\min} \leq \mu_{ij} < \mu_{ij}^{(\min)-1} \end{pmatrix}$$

$$(5)$$

Using (4) and (5), we can formulate the second constraint that satisfies budget limitation, in the following way:

$$\sum C_{ij}(\mu_{ij}) \le M \tag{6}$$

So, the proposed TCTP can be summarized as a nonlinear optimization model, as follows:

Max { Min
$$\varphi\left(\frac{T_d - \sum_{ij \in r} \mu_{ij}}{\sum_{ij \in r} \sigma_{ij}}\right) \quad \forall \ r=1,2,...,L \ }$$
 (7)

Subject to:

$$\sum C_{ij}(\mu_{ij}) \le M \,, \tag{8}$$

$$\mu_{ij}^{l} \le \mu_{ij} \le \mu_{ij}^{u} \quad \forall Activity \ ij \in r$$
(9)

This is a new formulation of TCTP in PERT networks. In order to shorten each activity, some amount of extra resources is needed. This amount of resources may be provided by external suppliers, who may determine some incentive policy such as price discount. Our proposed TCTP, models this situation. Using this model, project manager can increase confidence level of on-time project completion and prevent project delay. Also, proposed

model help us to use the available extra budget more efficiently, when we could find suppliers with discount policies.

3. Applying Genetic Algorithm

Genetic Algorithm (GA) is one of the popular metaheuristic algorithms, which can explore the feasible space using computational intelligence inspired from natural genetics and evolutionary concepts. It is a search method based on random selection policy can be used to solve the nonlinear mathematical programs. In this technique, an initial population containing several feasible solutions (chromosome) is generated randomly. Interesting fact is that GA does not need a good initial solution. In this procedure, algorithm starts from an initial solution and then it would be improved through an evolutionary process. After production of initial randomly generated population, parents are selected from this primary population and produce offspring. Second population is a combination of parent and offspring. The generation procedure has been done using two effective operators, cross over and mutation. The crossover is a genetic operator used to vary the programming of a chromosome or chromosomes from one generation to the next and the mutation operator can prevent the premature convergence of a new

generation [28]. Desirability of obtained offspring is investigated using fitness function which is concerned with objective function and feasibility. The most common steps for GA follow:

- Generate the initial population
- Evaluate the fitness of all solutions in population
- Repeat: Select parents

Apply the cross over and mutation operators to produce children

Evaluate the fitness of children

Establish new population using a combination of parents and children

• Until a satisfactory solution has been obtained About the proposed GA for the discount based TCTP, we develop the chromosome structure as a simple string of decision variables. Indeed, each chromosome represents μ_{ii} for all of the activities. Proposed GA algorithm

steps are presented in the following way:

- Initial random population is generated using uniform random number generator.
- Fitness of individuals is evaluated.
- Population is ranked in terms of computed fitness.
- Parents are selected randomly.
- Cross over and mutation are applied to produce children.
- Produced offspring are evaluated using fitness function.
- Form the new population using half of the fittest parent and half of the fittest children.

About the fitness function some points should be mentioned. In order to evaluate a solution (chromosome), two parameter should be investigated, goodness of solution and meaningfulness of solution. The goodness is concerned with objective function (project completion probability) and the meaningfulness indicates the feasibility. According to these parameters the fitness function can be represented as follows:

Fitness Function =
$$\prod_{ij \in project} \gamma_{ij} \times \left[Min \ \varphi \left(\frac{T_d - \sum_{ij \in r} \mu_{ij}}{\sum_{ij \in r} \sigma_{ij}} \right) \ \forall \ r = 1, 2, ..., L \right]$$
 (10)

In which, the first term represents feasibility and second one is concerned with project completion probability. γ_{ij} would be equal to one, if $\mu_{ij}^l \le \mu_{ij} \le \mu_{ij}^u \& \sum C_{ij}(\mu_{ij}) \le M$ and otherwise, indicates zero.

4. Numerical Example

In order to investigate the performance of the presented GA approach for the proposed discount based model of stochastic TCTP, a numerical example are described in this section. We conducted the proposed approach for the large scale example with 15 nodes, 20 independent activities and 44 interrelated paths to show that how the presented approach optimally improves the project completion probability. Characteristics of this network are presented in Table 2. The objective is to obtain optimal allocated budget to the activities in order to improve the all path completion probability from a risky value to a maximum confident one. It is assumed that the time unit is in weeks and the cost is in thousand dollars. According to the presented characteristics, initial probability of the project completion time at $T_d = 150$ which can be computed by using the Central Limit Theorem (CLT) is equal

to 55%. This value is related to a situation that all activities are planned in their upper bound of μ_{ij} . The proposed approach attempts to improve this value to a maximum level using the limited available budget in order to decrease the risk of project tardiness. It is also assumed that the available amount of additional budget for this purpose is equal to 10,000 thousand dollars. We conducted our analysis of the presented example in two steps. In first step we run the developed procedure with the random initial values for parameters of the proposed GA. Then the obtained results are considered to determine the efficient values of the parameters in terms of the accuracy and computation time. For this purpose we developed a computer simulation program based on the proposed approach and randomly initial values of GA parameters have been considered in primary computational effort. The obtained simulation results for the considered example are organized in Table 1.

To investigate efficiency of the proposed GA approach, best obtained project completion probability and CPU time have been considered. In order to make trade off between the accuracy and computation time, we also solve the example with the different values of population size (k), mutation and cross over probabilities, then

μ_{ij}^l	μ^1_{ij}	μ_{ij}°	μ^u_{ij}	S_{ij}^2	S_{ij}^1	S_{ij}°	σ_{ij}
8.51	12.32	13.29	14.26	201.59	250.00	265.95	0.25
11.60	-	14.56	16.70	-	195.25	225.45	0.60
9.56	-	10.05	15.09	-	302.70	321.28	0.45
12.95	15.29	16.09	17.77	259.30	295.52	305.04	0.15
11.72	12.56	14.05	16.79	252.27	265.33	294.37	0.51
11.14	13.00	15.25	16.34	257.84	301.04	326.65	0.67
10.43	-	12.05	15.77	-	254.05	295.73	0.65
13.43	14.59	17.05	18.15	194.51	201.36	254.00	0.41
13.03	15.05	16.32	17.83	241.08	259.32	294.21	0.62
12.61	-	15.02	17.50	-	201.51	285.54	0.57
9.46	10.25	12.98	15.01	199.65	209.21	225.31	0.14
13.22	-	16.02	17.99	-	295.84	309.74	0.03
12.41	-	17.02	17.34	-	229.74	276.48	0.24
11.32	-	15.25	16.48	-	207.97	257.54	0.10
10.17	14.52	15.02	15.57	154.95	174.04	198.74	0.58
8.58	10.21	14.09	14.32	174.00	198.32	254.13	0.36
9.00	-	12.06	14.65	-	205.27	211.37	0.50
10.50	11.65	14.21	15.84	298.57	314.45	365.36	0.25
9.31	-	14.09	14.89	-	302.12	341.91	0.17
	8.51 11.60 9.56 12.95 11.72 11.14 10.43 13.43 13.03 12.61 9.46 13.22 12.41 11.32 10.17 8.58 9.00 10.50	8.51 12.32 11.60 - 9.56 - 12.95 15.29 11.72 12.56 11.14 13.00 10.43 - 13.43 14.59 13.03 15.05 12.61 - 9.46 10.25 13.22 - 12.41 - 11.32 - 10.17 14.52 8.58 10.21 9.00 - 10.50 11.65	8.51 12.32 13.29 11.60 - 14.56 9.56 - 10.05 12.95 15.29 16.09 11.72 12.56 14.05 11.14 13.00 15.25 10.43 - 12.05 13.43 14.59 17.05 13.03 15.05 16.32 12.61 - 15.02 9.46 10.25 12.98 13.22 - 16.02 12.41 - 17.02 11.32 - 15.25 10.17 14.52 15.02 8.58 10.21 14.09 9.00 - 12.06 10.50 11.65 14.21	8.51 12.32 13.29 14.26 11.60 - 14.56 16.70 9.56 - 10.05 15.09 12.95 15.29 16.09 17.77 11.72 12.56 14.05 16.79 11.14 13.00 15.25 16.34 10.43 - 12.05 15.77 13.43 14.59 17.05 18.15 13.03 15.05 16.32 17.83 12.61 - 15.02 17.50 9.46 10.25 12.98 15.01 13.22 - 16.02 17.99 12.41 - 17.02 17.34 11.32 - 15.25 16.48 10.17 14.52 15.02 15.57 8.58 10.21 14.09 14.32 9.00 - 12.06 14.65 10.50 11.65 14.21 15.84	8.51 12.32 13.29 14.26 201.59 11.60 - 14.56 16.70 - 9.56 - 10.05 15.09 - 12.95 15.29 16.09 17.77 259.30 11.72 12.56 14.05 16.79 252.27 11.14 13.00 15.25 16.34 257.84 10.43 - 12.05 15.77 - 13.43 14.59 17.05 18.15 194.51 13.03 15.05 16.32 17.83 241.08 12.61 - 15.02 17.50 - 9.46 10.25 12.98 15.01 199.65 13.22 - 16.02 17.99 - 12.41 - 17.02 17.34 - 11.32 - 15.25 16.48 - 10.17 14.52 15.02 15.57 154.95 8.58 10.21 14.09 14.32 174.00 9.00 - 12.06 14.65 - </td <td>8.51 12.32 13.29 14.26 201.59 250.00 11.60 - 14.56 16.70 - 195.25 9.56 - 10.05 15.09 - 302.70 12.95 15.29 16.09 17.77 259.30 295.52 11.72 12.56 14.05 16.79 252.27 265.33 11.14 13.00 15.25 16.34 257.84 301.04 10.43 - 12.05 15.77 - 254.05 13.43 14.59 17.05 18.15 194.51 201.36 13.03 15.05 16.32 17.83 241.08 259.32 12.61 - 15.02 17.50 - 201.51 9.46 10.25 12.98 15.01 199.65 209.21 13.22 - 16.02 17.99 - 295.84 12.41 - 17.02 17.34 - 229.74 11.32 - 15.25 16.48 - 207.97 10.17 14</td> <td>8.51 12.32 13.29 14.26 201.59 250.00 265.95 11.60 - 14.56 16.70 - 195.25 225.45 9.56 - 10.05 15.09 - 302.70 321.28 12.95 15.29 16.09 17.77 259.30 295.52 305.04 11.72 12.56 14.05 16.79 252.27 265.33 294.37 11.14 13.00 15.25 16.34 257.84 301.04 326.65 10.43 - 12.05 15.77 - 254.05 295.73 13.43 14.59 17.05 18.15 194.51 201.36 254.00 13.03 15.05 16.32 17.83 241.08 259.32 294.21 12.61 - 15.02 17.50 - 201.51 285.54 9.46 10.25 12.98 15.01 199.65 209.21 225.31 13.22 - 16.02 17.99 - 295.84 309.74 12.41 - <</td>	8.51 12.32 13.29 14.26 201.59 250.00 11.60 - 14.56 16.70 - 195.25 9.56 - 10.05 15.09 - 302.70 12.95 15.29 16.09 17.77 259.30 295.52 11.72 12.56 14.05 16.79 252.27 265.33 11.14 13.00 15.25 16.34 257.84 301.04 10.43 - 12.05 15.77 - 254.05 13.43 14.59 17.05 18.15 194.51 201.36 13.03 15.05 16.32 17.83 241.08 259.32 12.61 - 15.02 17.50 - 201.51 9.46 10.25 12.98 15.01 199.65 209.21 13.22 - 16.02 17.99 - 295.84 12.41 - 17.02 17.34 - 229.74 11.32 - 15.25 16.48 - 207.97 10.17 14	8.51 12.32 13.29 14.26 201.59 250.00 265.95 11.60 - 14.56 16.70 - 195.25 225.45 9.56 - 10.05 15.09 - 302.70 321.28 12.95 15.29 16.09 17.77 259.30 295.52 305.04 11.72 12.56 14.05 16.79 252.27 265.33 294.37 11.14 13.00 15.25 16.34 257.84 301.04 326.65 10.43 - 12.05 15.77 - 254.05 295.73 13.43 14.59 17.05 18.15 194.51 201.36 254.00 13.03 15.05 16.32 17.83 241.08 259.32 294.21 12.61 - 15.02 17.50 - 201.51 285.54 9.46 10.25 12.98 15.01 199.65 209.21 225.31 13.22 - 16.02 17.99 - 295.84 309.74 12.41 - <

14.97

Table 1. Characteristics of considered example.

compute the computation time and efficiency measure. The appropriate levels of these parameters can conduct the GA toward optimal solution, directly and therefore, the analysis of these parameters is an important aspect of the proposed GA approach. Tables 2, 3 and 4 provide the behavior of the model results according to different levels of parameters, for presented example.

11.64

13.92

9.41

According to the simulation results provided by Tables 2, 3 and 4 the best probability of cross over and mutation rates are 0.8 and 0.2, respectively and best values of considered criteria (accuracy and CPU time) obtained by setting population size at 25-30. In this example the project completion probability has been improved from 55% to 0.79%.

5. Model Validation

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The appropriate levels of the GA parameters have been investigated in section 4 to reach the accurate results with reasonable computation time. By using this procedure we reach the maximum capabilities of the proposed GA approach in terms of accuracy and computational time, but these maximum capabilities should be compared

Table 2. Model results (objective function) for different levels of P_C ($P_m=0.2$, ${\bf k}=$ 25).

164.18

187.47

0.63

P_c	Run 1*	Run 2	Run 3	Run 4	Best Obj.
0.1	0.7390	0.7142	0.6952	0.7202	0.7390
0.2	0.6966	0.6414	0.7157	0.7377	0.7377
0.4	0.7222	0.6197	0.7585	0.6726	0.7585
0.6	0.7247	0.6253	0.7822	0.6792	0.7822
0.8	0.7009	0.6957	0.7951	0.7823	0.7951

Each runs containing 200 iterations

157.37

Table 3. Model results (objective function) for different levels of P_m ($P_c=0.7$, ${\bf k}={\bf 25}$).

$P_{\scriptscriptstyle m}$	Run 1	Run 2	Run 3	Run 4	Best Obj.
0.1	0.7108	0.7561	0.6970	0.7896	0.7896
0.2	0.6876	0.7857	0.7952	0.7396	0.7952
0.4	0.7185	0.6969	0.7355	0.7291	0.7355
0.6	0.6827	0.6873	0.6269	0.7565	0.7565
0.8	0.7007	0.7067	0.7477	0.6900	0.7477

Run 1 Run 2 Run 3 Run 4 Best k **CPU Time** Obj. CPU CPU CPU CPU Obj. Obj. Func. Obj. Func. Obj. Func. Time Time time Func. Time 0.701 5 10.40 0.692 11.24 0.642 9.80 0.716 15.59 10.40 0.716 0.707 10 16.10 0.722 15.92 25.21 0.757 20.65 0.740 15.92 0.757 12 17.41 0.710 16.10 0.735 29.51 0.740 23.62 0.727 16.10 0.740 18 30.60 0.731 29.25 0.765 39.74 0.788 45.41 0.751 29.25 0.788 25 40.25 0.758 45.62 0.761 54.30 0.780 42.20 0.796 40.25 0.796 51.50 0.793 47.26 0.757 59.20 0.798 65.20 0.795 47.26 0.798 30

Table 4. Model results (computational time, objective function) for different levels of k ($P_C = 0.7$, $P_m = 0.2$).

with other standard solutions which are available in the literature to validate the solution. Since, this paper is the first combination of price discount with the TCTP, there isn't any standard test problem for the proposed model, in the published literature. We use the global optimum solution obtained by LINGO software to evaluate the performance of proposed approach. Indeed, LINGO can handle nonlinear programming problems involving both continuous and binary variables and solve such problem by generalized reduced gradient (GRG).

For this purpose a mathematical model based on a single path of PERT network has been considered. The results of comparing the proposed GA's best objective function (obtained by appropriate levels of GA parameters based on the procedure described in section 4) and

LINGO's global optimum for a single path are presented in Table 5.

According to the Table 5, the proposed GA approach shows a difference with LINGO's global optimum with the average of 1.84%. Such a little difference can be reliable and shows the good performance of proposed approach.

6. Conclusions

In this paper, we proposed a new approach based on price discount for TCTP in PERT networks in which activity durations follow the normal distribution. The main objective of the developed model is to improve the

Table 5. Comparing the proposed GA best objective function and LINGO's global optimum.

Problem	Limited Budget (M)	Proposed GA	LINGO	Differences (LINGO-GA)	% Differences
1	10,000	0.7425	0.7489	0.0064	0.8620
2	15,000	0.7689	0.7909	0.0220	2.8612
3	20,000	0.7849	0.7981	0.0132	1.6817
4	25,000	0.8001	0.8214	0.0213	2.6622
5	30,000	0.8149	0.8236	0.0087	1.0676
6	35,000	0.8311	0.8544	0.0233	2.8035
7	40,000	0.864	0.8764	0.0124	1.4352
8	60,000	0.9049	0.9171	0.0122	1.3482
Average					1.8402 %

project completion probability from a risky amount to a maximum value using limited additional budget. To our knowledge, this is the first combination of price discount as a supplier's policy, with the TCTP, in the published literature. The presented model determines the optimal additional resources should be allocated to activities using the genetic algorithm. The GA algorithm has organized to allocate the available budget to activities. In order to illustrate the model efficiency, a computer program using MATLAB 7.6.0 was developed and the model was tested on a medium scale PERT network. Best amount of objective function and CPU time have been recorded and efficient levels of GA parameters have been investigated through the several computational experiments. In order to validate the GA approach, proposed model have been compared with the LINGO's global optimum solution and obtained results showed the good performance of the proposed solution approach.

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